

# Strategic thinking in infinite games

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Museu de la Ciència CosmoCaixa  
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# Recursive chess



Etching by Django Pinter, Oxford.

# The Chocolatier's game



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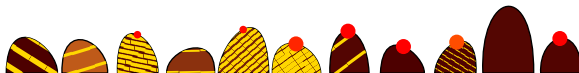


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# The Chocolatier's game



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Each round, the *Glutton* selects one and eats it.

Glutton wins, after infinite play, if all chocolates are consumed.

# Can the Glutton win?

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Chocolates steadily accumulate on the platter, and the Glutton falls further and further behind.

Nevertheless, with strategic choices, I claim that the Glutton CAN win. He can eat all the chocolate.



# How the Glutton wins

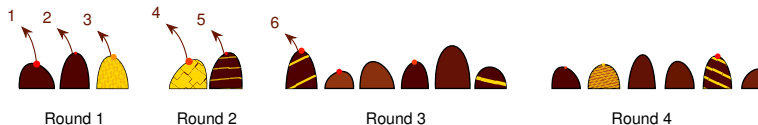
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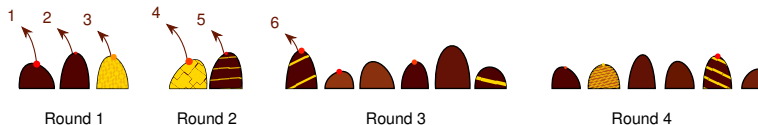


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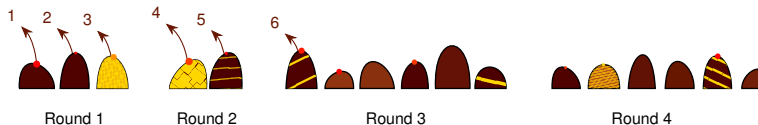
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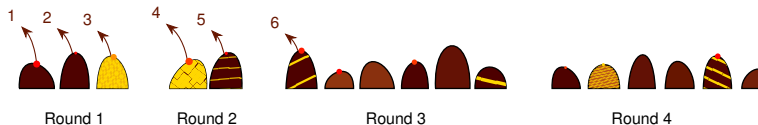
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Newly served chocolates are placed at the back of the queue.

But the Glutton eats always from the front.

Every chocolate is finite distance from the front.

So after infinitely many turns, every chocolate will be eaten!

# Infinite chocolates served each round

What if the Chocolatier serves up not merely finitely many chocolates each round, but a countable infinity of new chocolates each round?



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Nevertheless, the Glutton can still win!



# Winning with infinitely many chocolates each round

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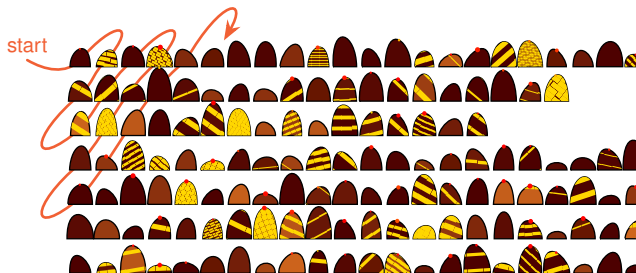
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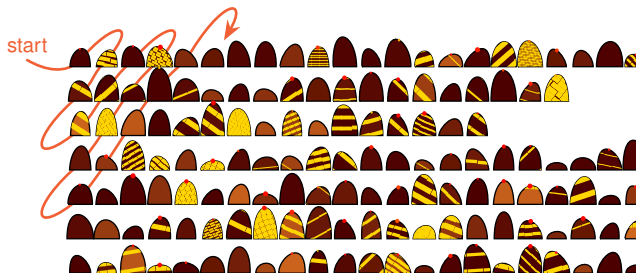


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Further considerations lead to deeper mathematical issues.

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- What if the Chocolatier is uncountably creative?
- In this case, there can be no memory-free winning strategy for the Glutton.
- Nevertheless, assuming the axiom of choice, the Glutton can win while knowing only chocolates on offer + the previously eaten chocolate.

# The Sudoku game

5			4			1		
3					5			7
	9				3	5		
2			7					
		4			8			
6								9
		6				4		
		1			9	2		
4				5			8	

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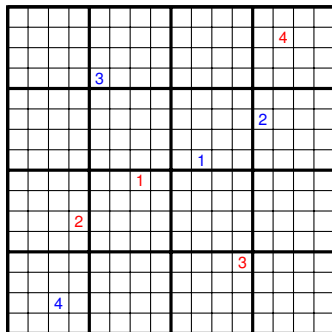
- Consider Sudoku as a two-player game
- Start from empty board, perhaps very large
- Players take turns placing numbers
- Plays must obey the Sudoku condition—no repeats on any row, column, or subboard
- A player loses when there is no legal move
- So this version of the game is not about a global solution

# Sudoku game, even size board

Second player has a winning strategy.

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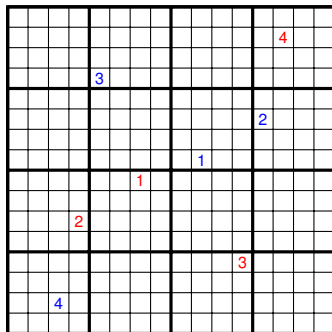
Strategy: Second player copies, reflected through the center.

Key observation: if previous move was legal, so is reflection.



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So the second player wins.

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So first player will win.

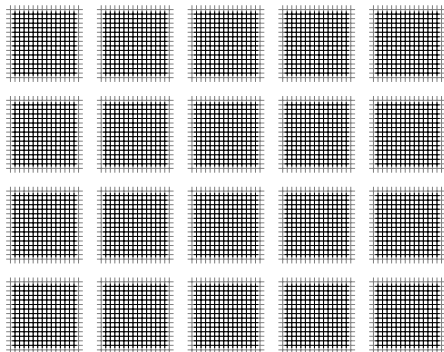
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But of course, I shall want to play *infinite* Sudoku!



# Infinite Sudoku

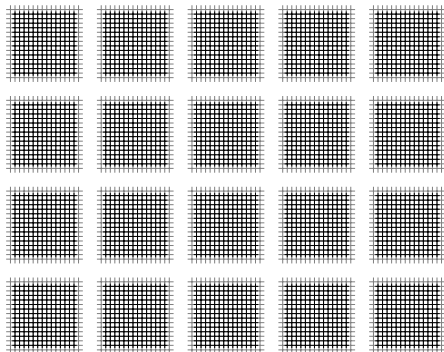
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Each subboard is like the integer grid.

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- The Spoiler aims to prevent this.

Can the Spoiler set traps and create obstacles so as to prevent the Solver from succeeding in the limit?

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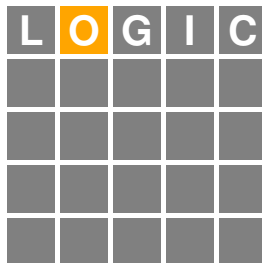
The Spoiler will be unable to prevent Solver from creating an infinite global solution.

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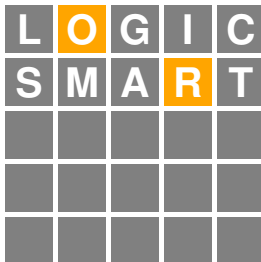
Solver can organize the requirements so as to fulfill them all.

Thus, the Solver has a winning strategy.

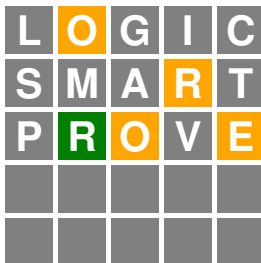
# Wordle



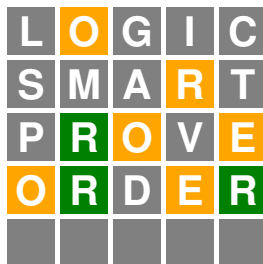
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S	U	P	E	R	C	A	L	I	F	R	A	G	I	S	T	I	C	E	X	P	I	A	L	I	D	O	C	I	O	U	S
A	N	T	I	D	I	S	E	S	T	A	B	L	I	S	H	M	E	N	T	A	R	I	A	N	I	S	M		1	2	3
I	T		W	A	S		A		D	A	R	K		A	N	D		S	T	O	R	M	Y		N	I	G	H	T	.	
O	N	E		F	I	S	H		T	W	O		F	I	S	H		R	E	D		F	I	S	H		B	L	U	E	

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Furthermore, let us consider any crazy fixed dictionary of allowed phrases, perhaps including many nonsense sequences of letters.

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With 26 letters, there are  $26^{1000000}$  many distinct “words” of length one million.

If we wrote down one million of them every second since the beginning of time at the Big Bang, we will have hardly started!

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Iterate. Since 26 letters, this can go on at most 26 steps.  $\square$

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Related to Infinite Mastermind. Things get very interesting and complicated.

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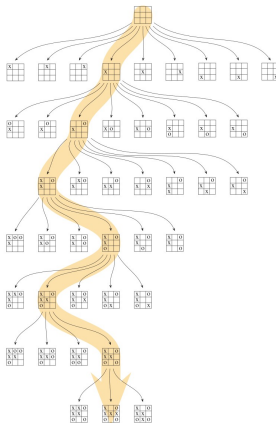
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- Does every infinite game have a winning strategy?

Some of these questions begin to engage with deep mathematical and sometimes philosophical questions.

# Game tree

Every game gives rise to a *game tree*, the tree of all possible positions of the game. Plays of the game are paths through the game tree.



# Fundamental theorem of finite games

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In every finite two-player game of perfect information, one of the players has a winning strategy.

I know five distinct proofs of this beautiful, classical theorem.

When teaching my logic classes at the university, I would usually spend several lectures giving all five proofs.

Let me describe just one of them here.

# Fundamental theorem—proof by back-propagation

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Working from bottom, label each node whether a player can win from that node.

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Whoever gets their label on the root will have a winning strategy—stay on your own label.  $\square$

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In any finite game of perfect information, allowing draws, either one of the players has a winning strategy or both players have draw-or-better strategies.

Follows from the win/loss version.

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## Chess

For the game of chess, either one of the players white or black has a winning strategy or both players have draw-or-better strategies.

# Open determinacy

Similar ideas work for *open* games, which are games for which all winning plays for one of the players are known as wins at a finite stage of play.

## Open determinacy theorem (Gale-Stewart 1953)

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## Open determinacy theorem (Gale-Stewart 1953)

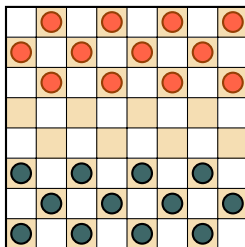
In every open game, one of the players has a winning strategy (or both have drawing strategies).

One way to prove this uses the concept of *transfinite ordinal game values*.

These generalize the chess idea of mate-in-2 or mate-in-3, but work with infinite games.

# Finite draughts

The familiar game of draughts, also known as checkers.

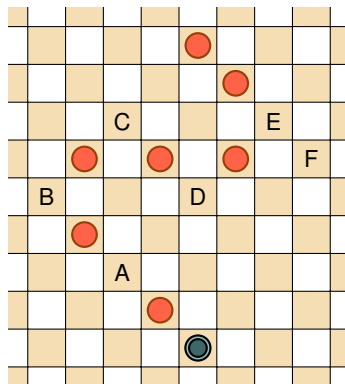


Although often played by children, draughts admits serious advanced play, with draughts grandmasters competing in international tournaments.



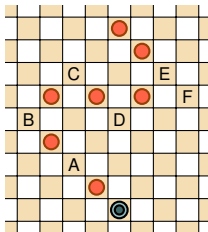
# Infinite draughts

Let us endeavor instead to play *infinite* draughts.



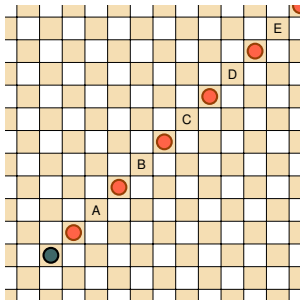
The checkerboard extends without end in every direction.

# The rules



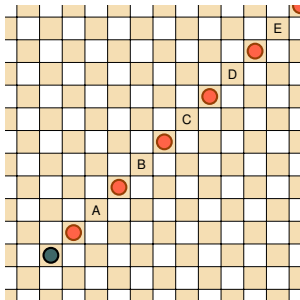
- No standard starting configuration.
- Analyze the game from any given position.
- Pawns and kings.
- Obligatory jumping rule.
- Obligatory iterated jumping.
- Winning condition: you lose when you have no legal move.

# Infinite iterations



During an infinite iterated jump, the red pieces are all removed.

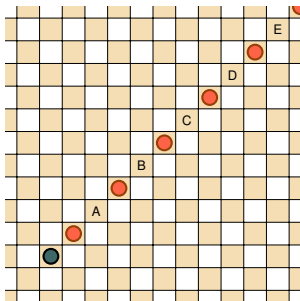
# Infinite iterations



During an infinite iterated jump, the red pieces are all removed.

A conundrum: what about the black piece?

# Infinite iterations

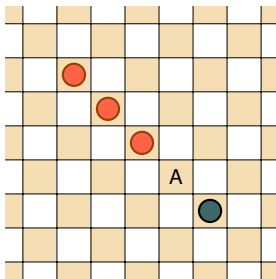


During an infinite iterated jump, the red pieces are all removed.

A conundrum: what about the black piece?

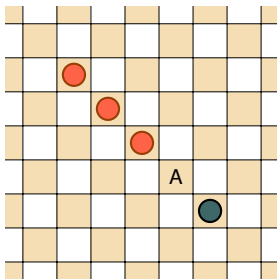
Infinite iterated-jump rule: the jumping piece also is removed.

## Game value 2



This position, with Red to play, has game value 2 for Red.

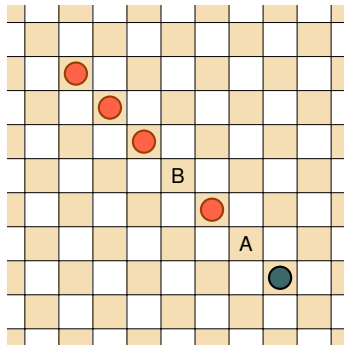
## Game value 2



This position, with Red to play, has game value 2 for Red.

Red can advance leading pawn to A as bait, obligating Black to jump, after which Red recaptures.

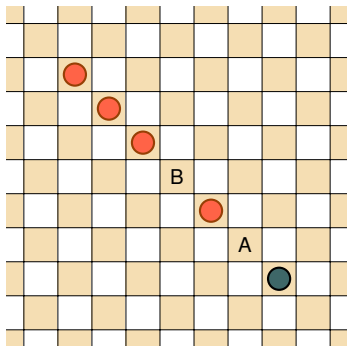
# Game value 3



Red to play, with game value 3.



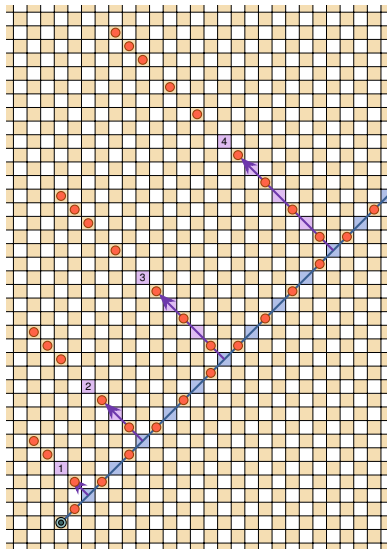
## Game value 3



Red to play, with game value 3.

Red advances isolated pawn to A, obligating Black to jump, then advance to B, Black jumps, then Red captures.

# Game value $\omega$



Black is obligated to jump.

Can jump infinitely,  
but that is an immediate loss.

So will come  
to rest at some square  $n$ .

And then lose in  $n$  moves.

This is game value  $\omega$ .

# Infinite draughts exhibits high game values

Theorem (Hamkins & Leonessi 2022 )

*Every countable ordinal arises as the game value of a position in infinite draughts.*

This result is optimal—infinite draughts exhibits the most robust spectrum of game values that is possible in a countable game.

# Infinite draughts exhibits high game values

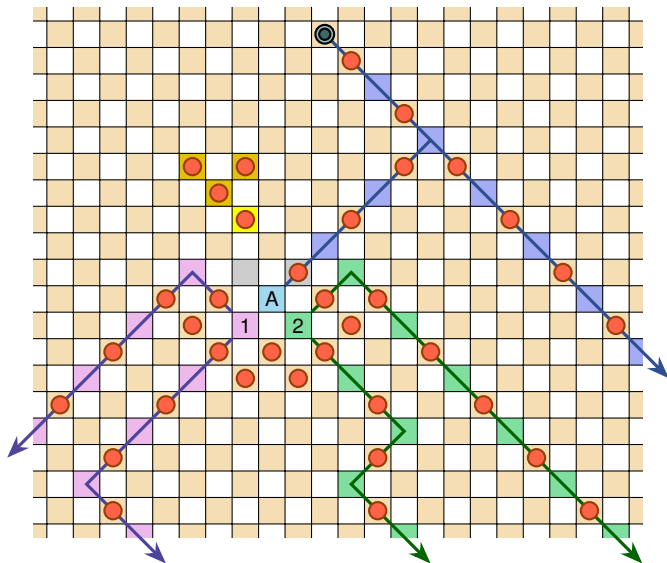
## Theorem (Hamkins & Leonessi 2022 )

*Every countable ordinal arises as the game value of a position in infinite draughts.*

This result is optimal—infinite draughts exhibits the most robust spectrum of game values that is possible in a countable game.

Main proof idea: embed certain well-founded trees into positions of infinite draughts.

The draughts play will unfold in a manner as though Black is climbing the tree, and consequently the game value will track the ordinal rank of the well-founded tree itself.



# Infinite chess

Let us turn now to infinite chess.

# Infinite chess

Let us turn now to infinite chess.

Infinite chess is a game of the mind—we do not sit down in a café and play infinite chess over espresso.

# Infinite chess

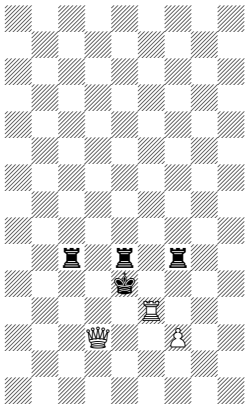
Let us turn now to infinite chess.

Infinite chess is a game of the mind—we do not sit down in a café and play infinite chess over espresso.

But rather, we sit down in that café and wonder what it would be like to play from this position or that one.

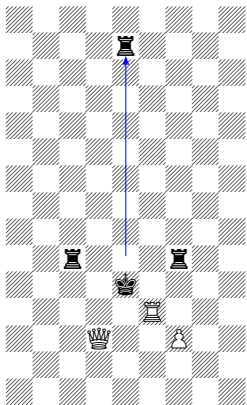


# A finite position with value $\omega$



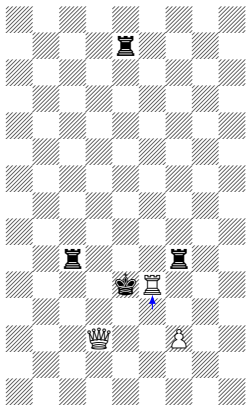
Black to move.

# A finite position with value $\omega$



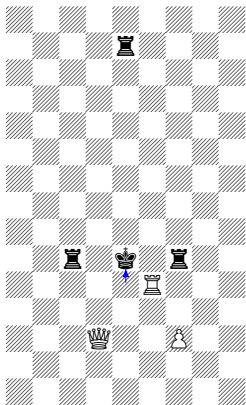
Black moves up arbitrary height

# A finite position with value $\omega$

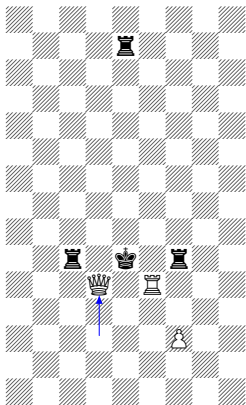


Check

# A finite position with value $\omega$

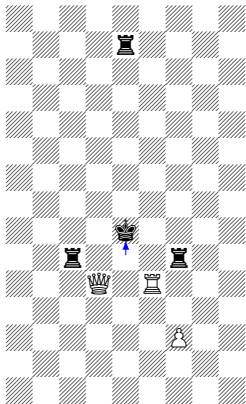


# A finite position with value $\omega$

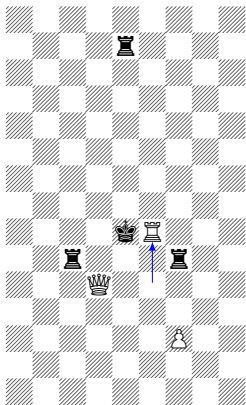


Check

# A finite position with value $\omega$



# A finite position with value $\omega$

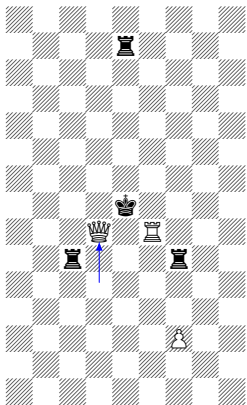


Check



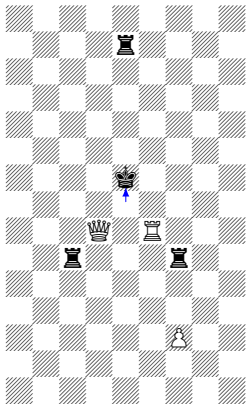


# A finite position with value $\omega$

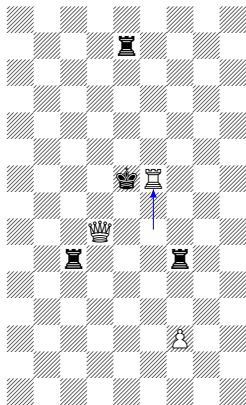


Check

# A finite position with value $\omega$

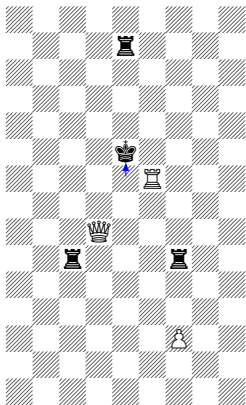


# A finite position with value $\omega$

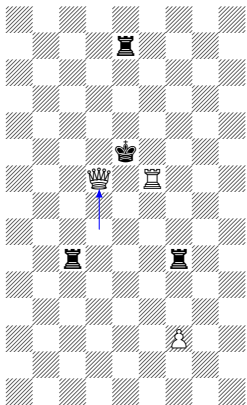


Check

# A finite position with value $\omega$

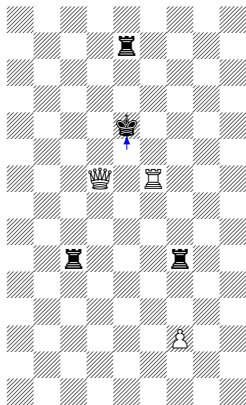


# A finite position with value $\omega$

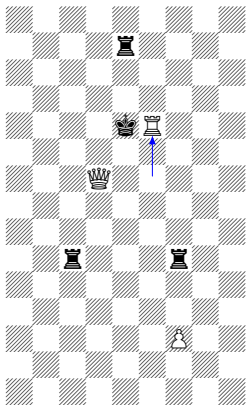


Check

# A finite position with value $\omega$

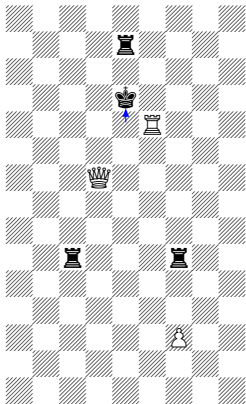


# A finite position with value $\omega$



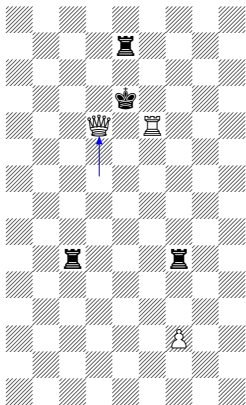
Check

# A finite position with value $\omega$



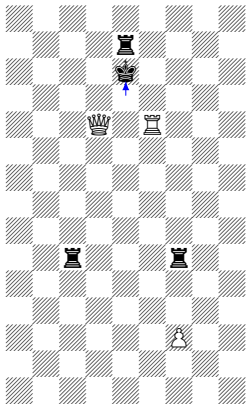


# A finite position with value $\omega$

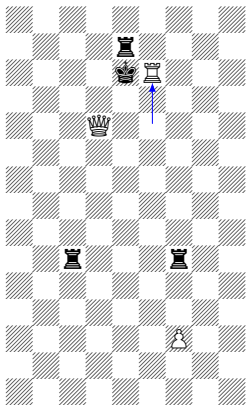


Check

# A finite position with value $\omega$

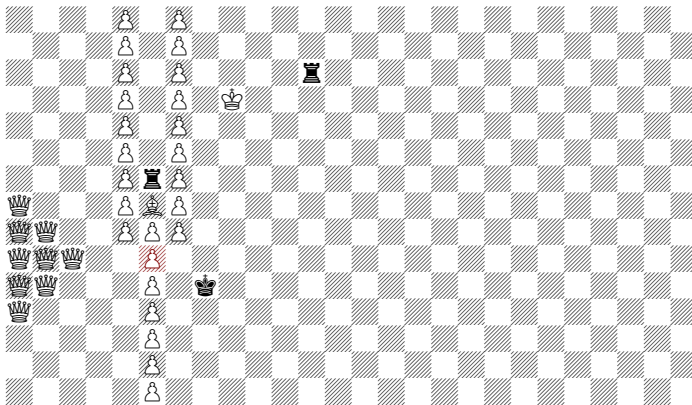


# A finite position with value $\omega$



Checkmate. Black can cause arbitrary delay, but the only choice is on first move, so the initial value is  $\omega$ .

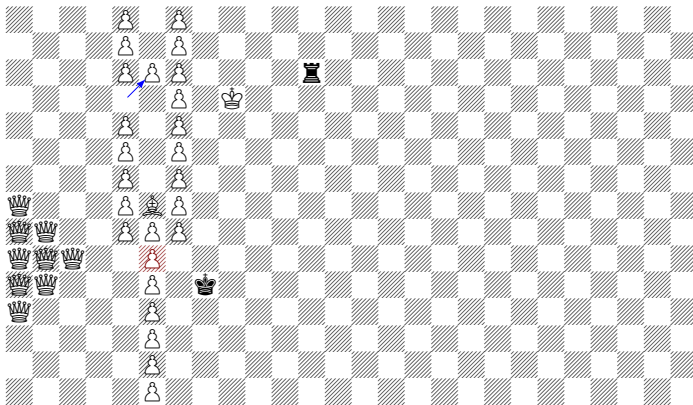
# Releasing the Hordes, with value $\omega^2$



Black to move.

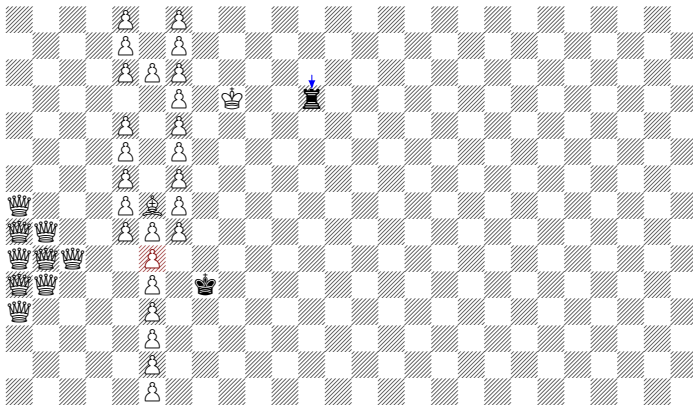
Joel David Hamkins

## Releasing the Hordes, with value $\omega^2$



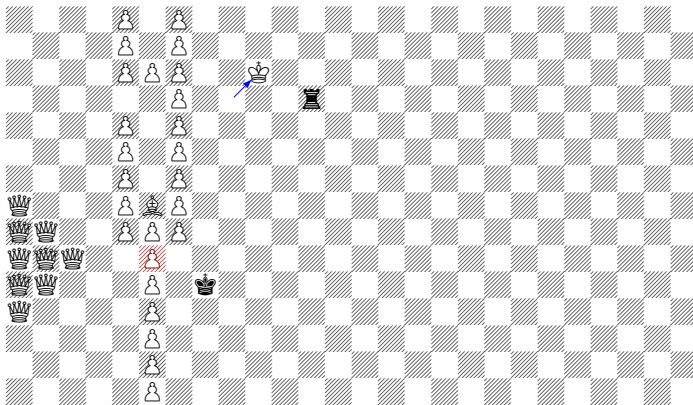
White should capture from left side.

# Releasing the Hordes, with value $\omega^2$



Now black begins to harass white king.

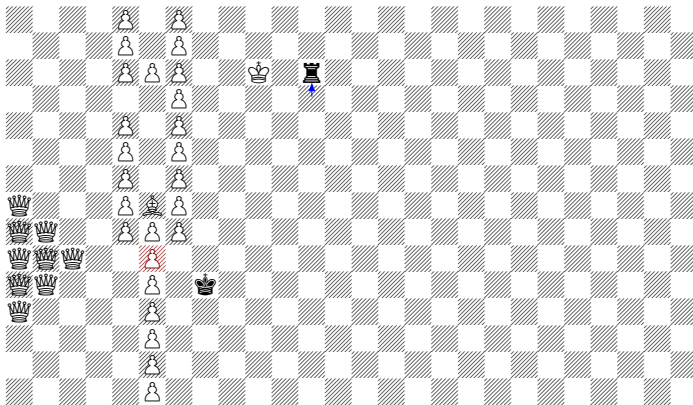
# Releasing the Hordes, with value $\omega^2$



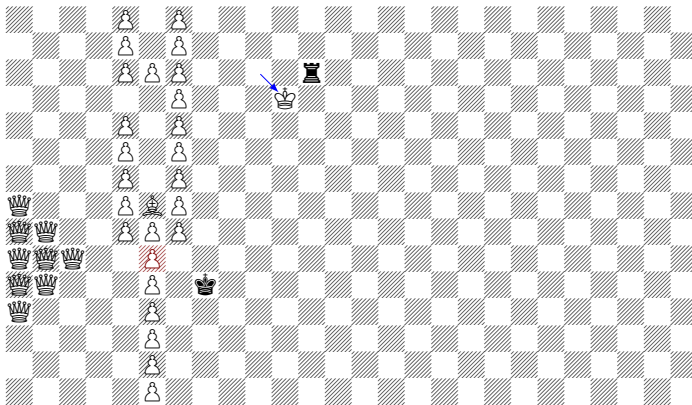
White must chase down the rook to avoid perpetual check.



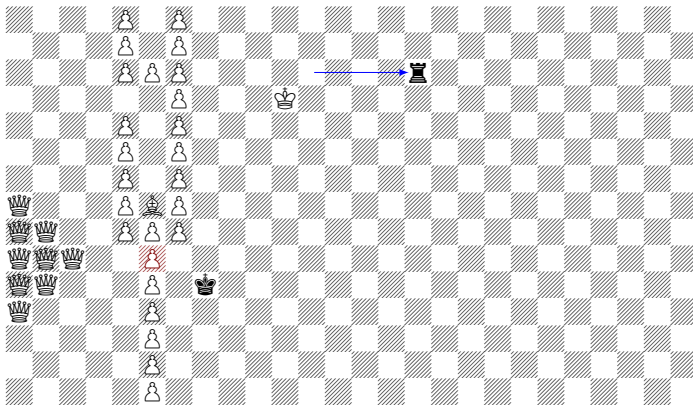
# Releasing the Hordes, with value $\omega^2$



## Releasing the Hordes, with value $\omega^2$

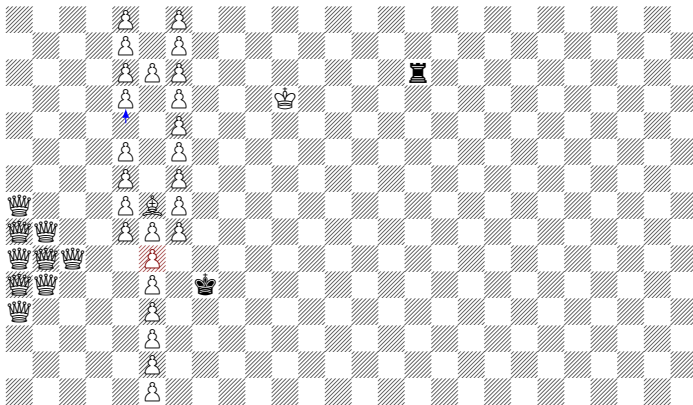


# Releasing the Hordes, with value $\omega^2$



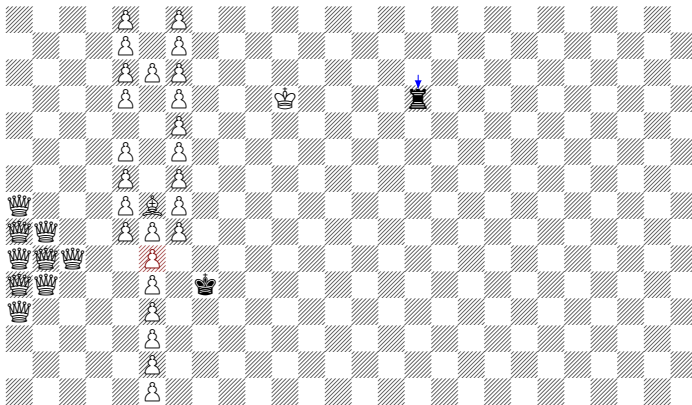
Black must move away to save rook.

# Releasing the Hordes, with value $\omega^2$

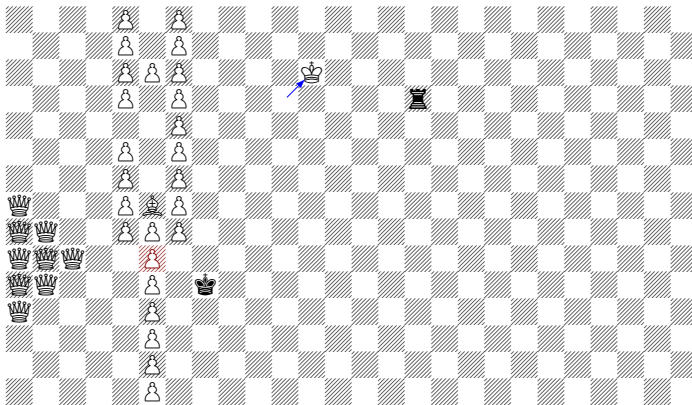


Now is white's chance to advance a pawn.

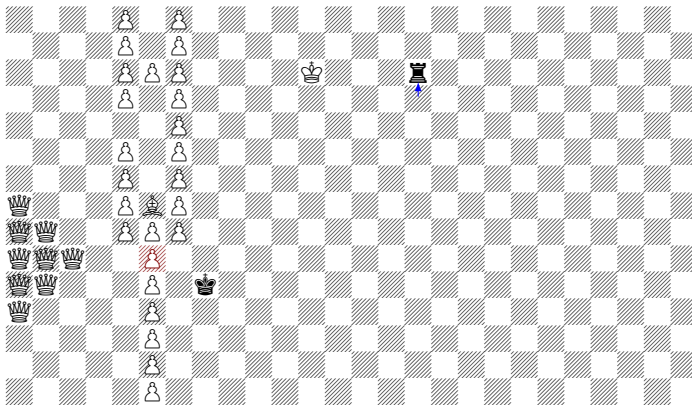
## Releasing the Hordes, with value $\omega^2$



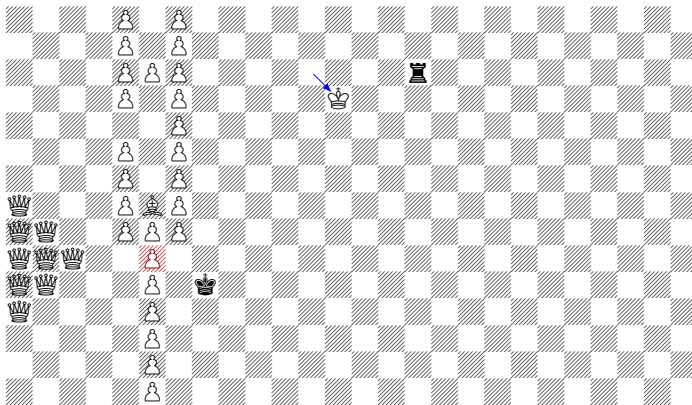
# Releasing the Hordes, with value $\omega^2$



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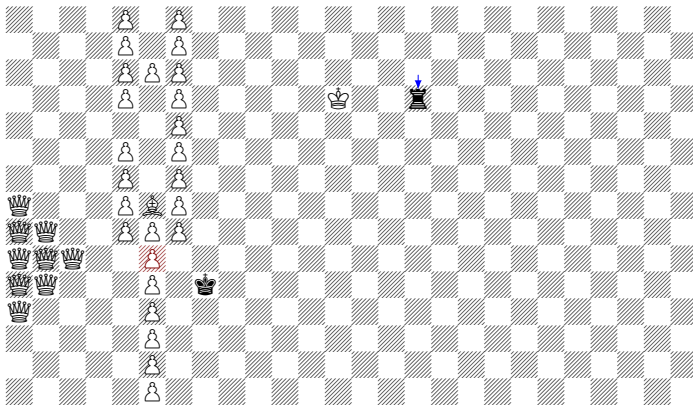


# Releasing the Hordes, with value $\omega^2$

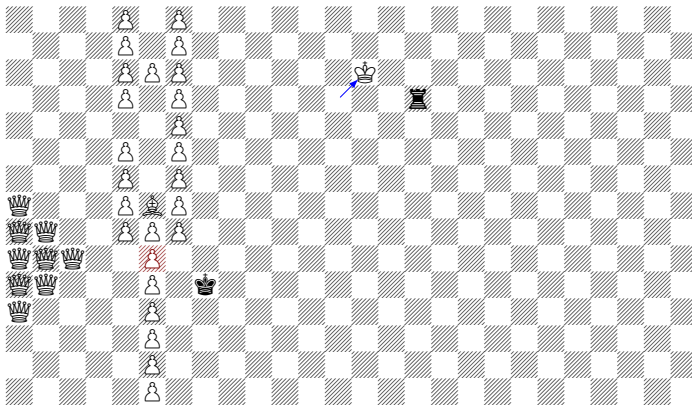




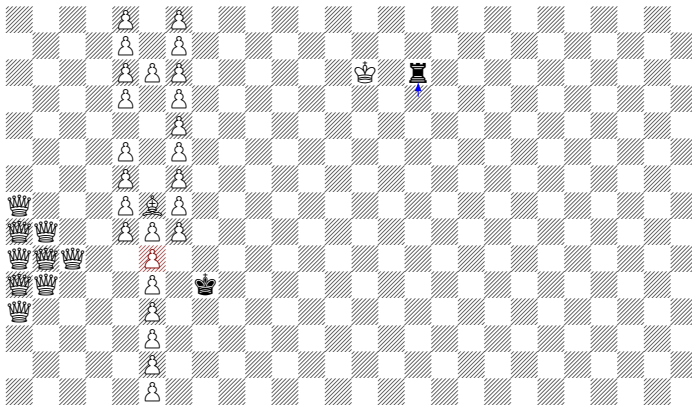
# Releasing the Hordes, with value $\omega^2$



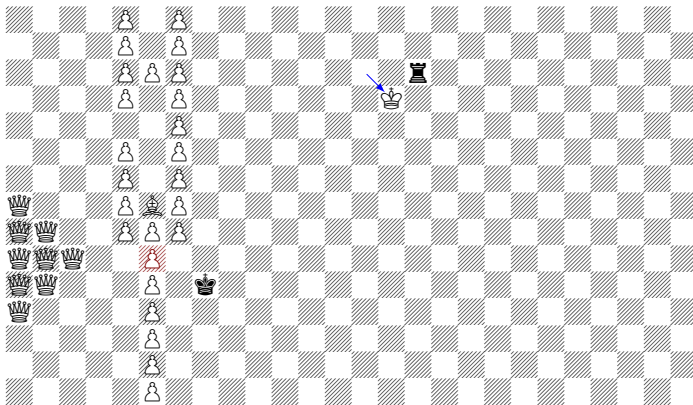
# Releasing the Hordes, with value $\omega^2$



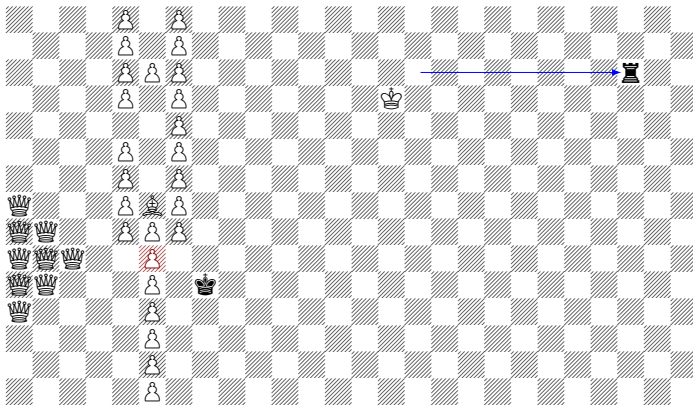
## Releasing the Hordes, with value $\omega^2$



# Releasing the Hordes, with value $\omega^2$

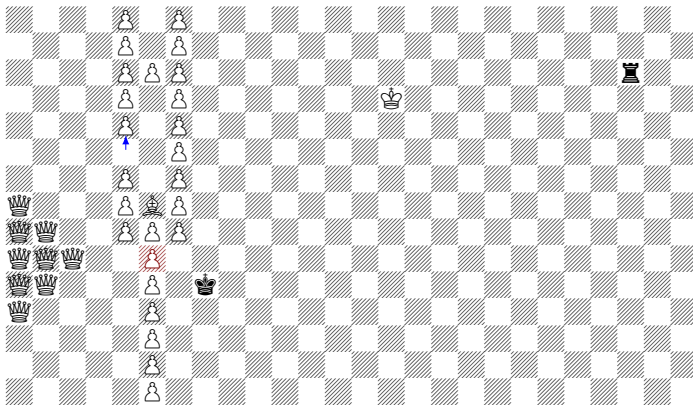


# Releasing the Hordes, with value $\omega^2$



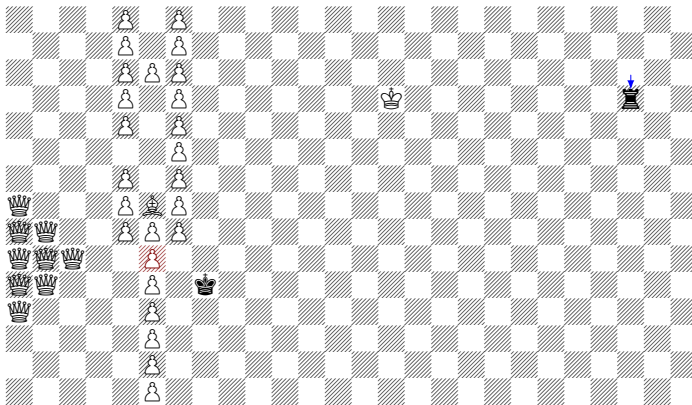
Black moves arbitrary distance out.

## Releasing the Hordes, with value $\omega^2$



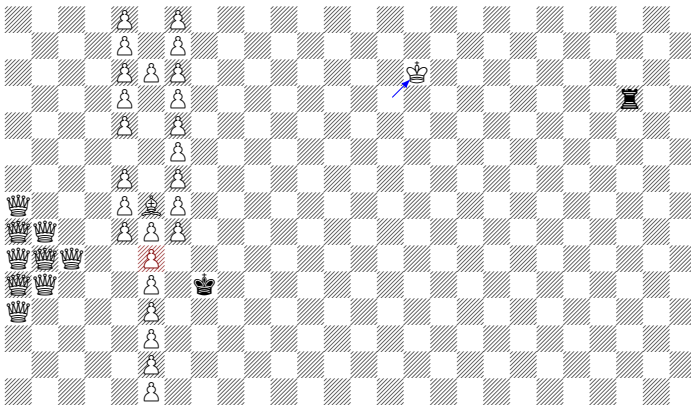
Another chance to advance a pawn.

# Releasing the Hordes, with value $\omega^2$



Black harasses the white king.

## Releasing the Hordes, with value $\omega^2$



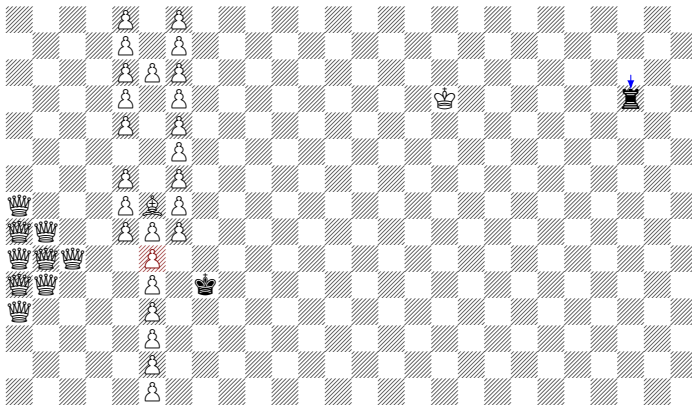
White must chase him down.



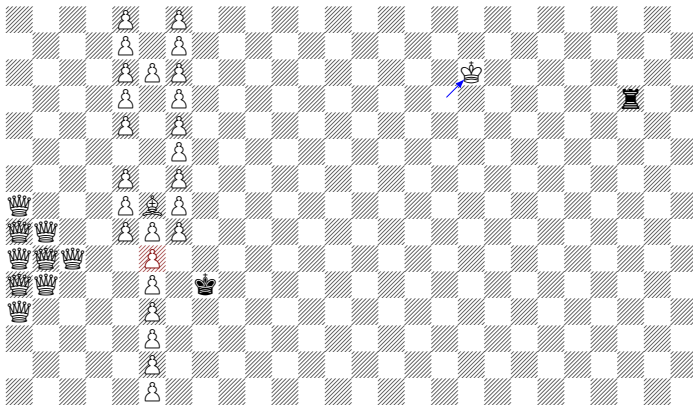




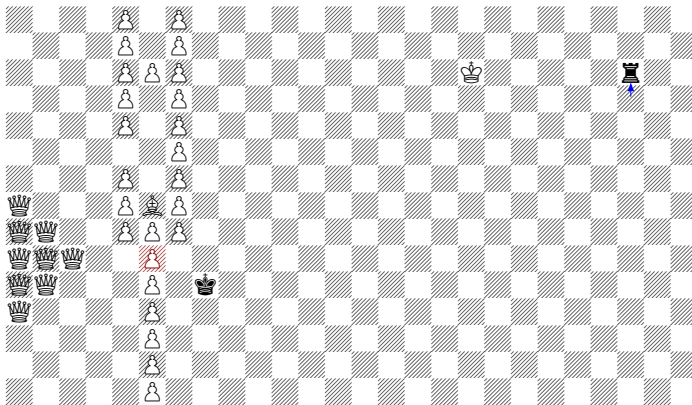
# Releasing the Hordes, with value $\omega^2$



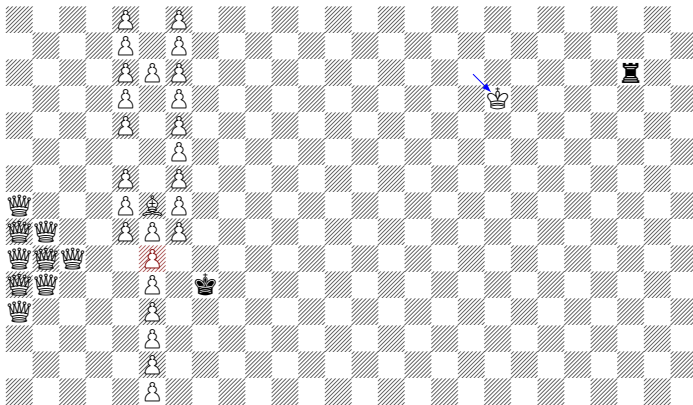
# Releasing the Hordes, with value $\omega^2$



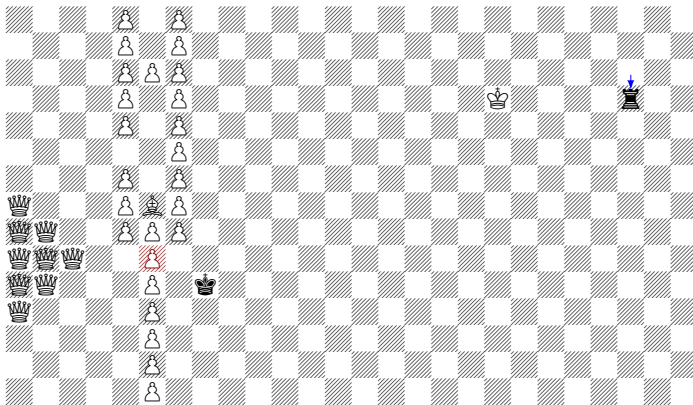
# Releasing the Hordes, with value $\omega^2$



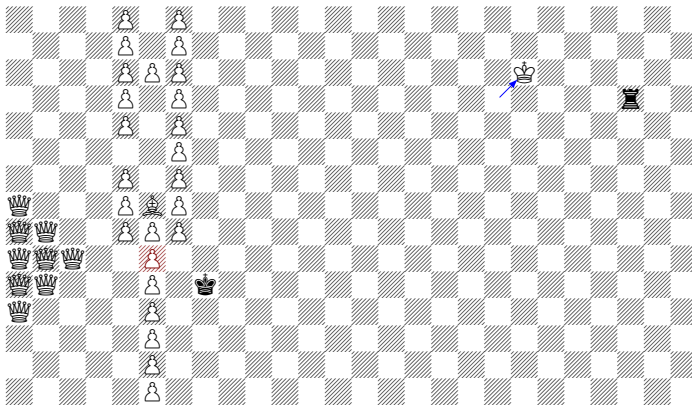
# Releasing the Hordes, with value $\omega^2$



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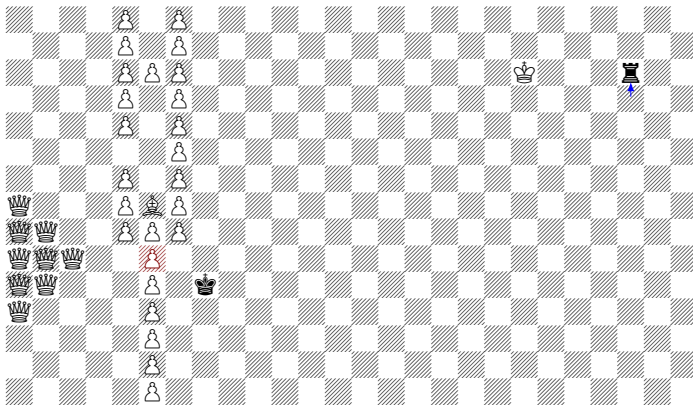


# Releasing the Hordes, with value $\omega^2$

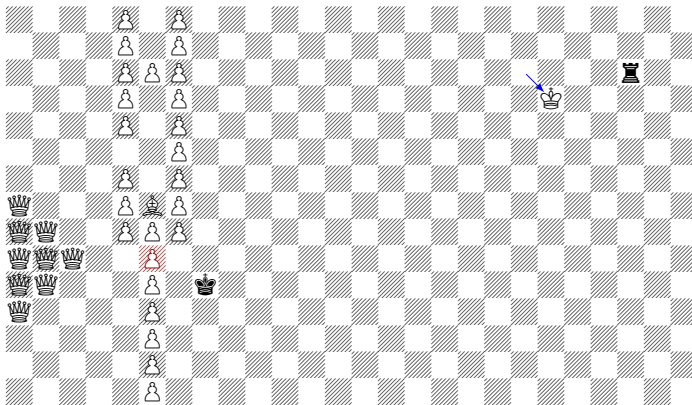




# Releasing the Hordes, with value $\omega^2$



## Releasing the Hordes, with value $\omega^2$



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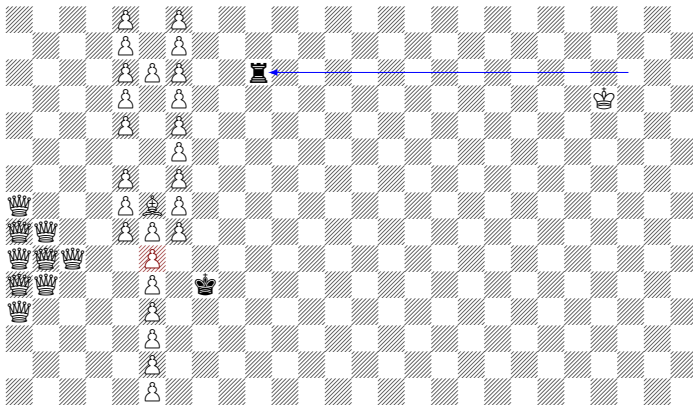
# Releasing the Hordes, with value $\omega^2$



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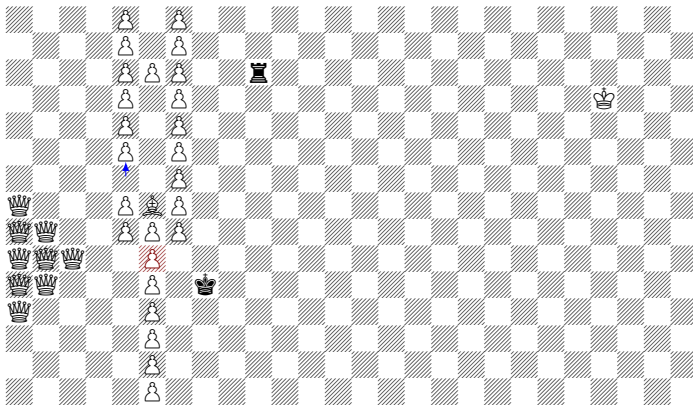


# Releasing the Hordes, with value $\omega^2$



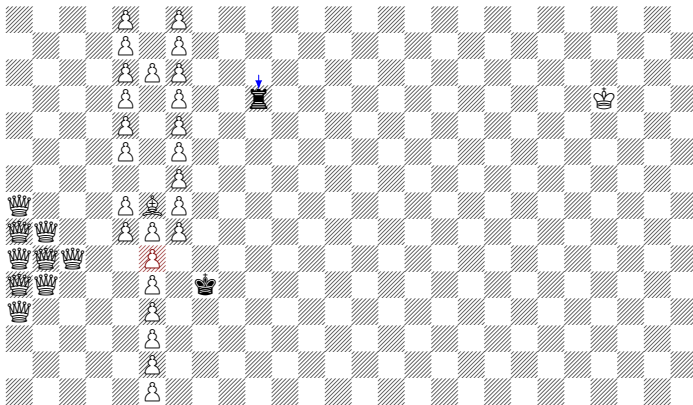
(Black should actually move arbitrary distance to the right.)

# Releasing the Hordes, with value $\omega^2$

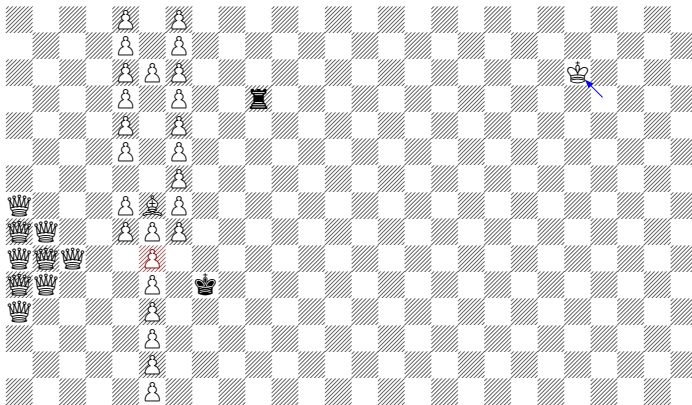




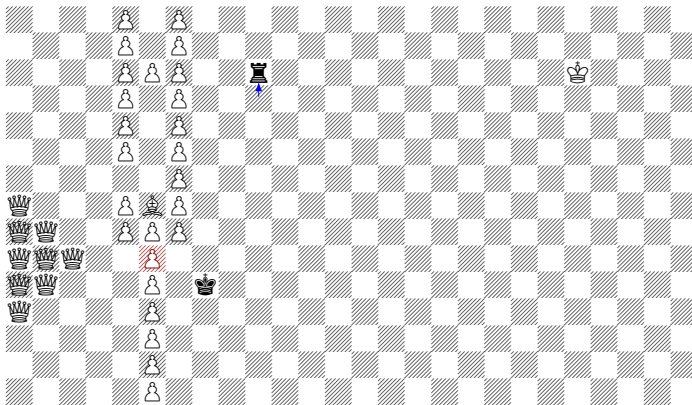
# Releasing the Hordes, with value $\omega^2$



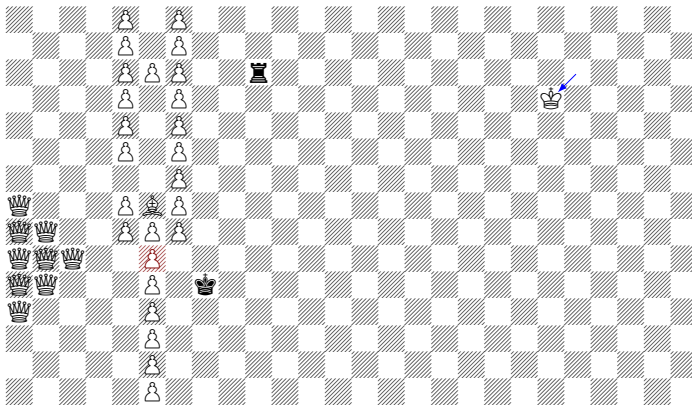
## Releasing the Hordes, with value $\omega^2$



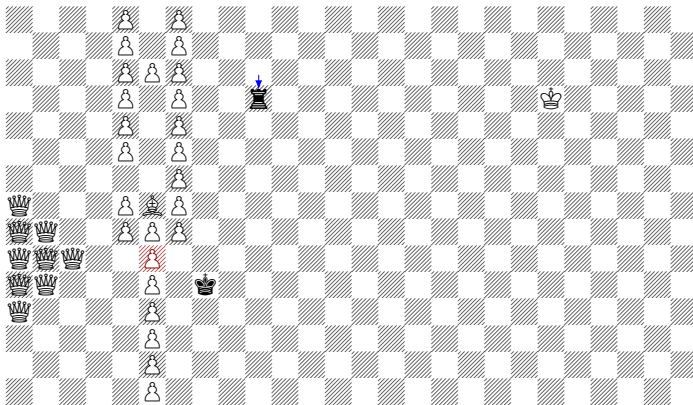
# Releasing the Hordes, with value $\omega^2$



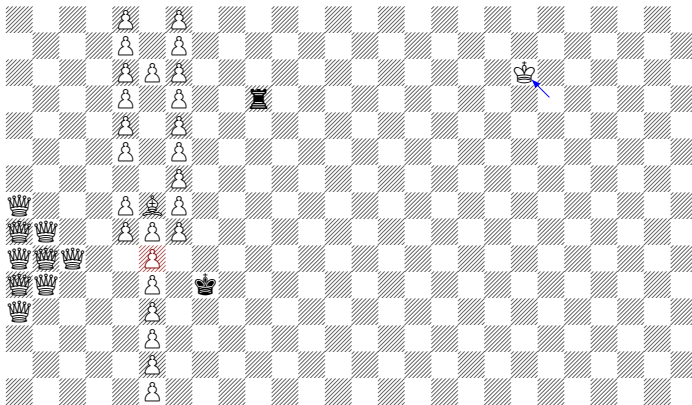
# Releasing the Hordes, with value $\omega^2$



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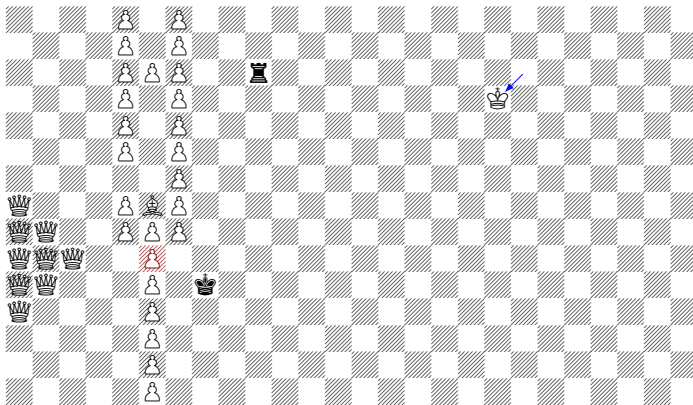


# Releasing the Hordes, with value $\omega^2$





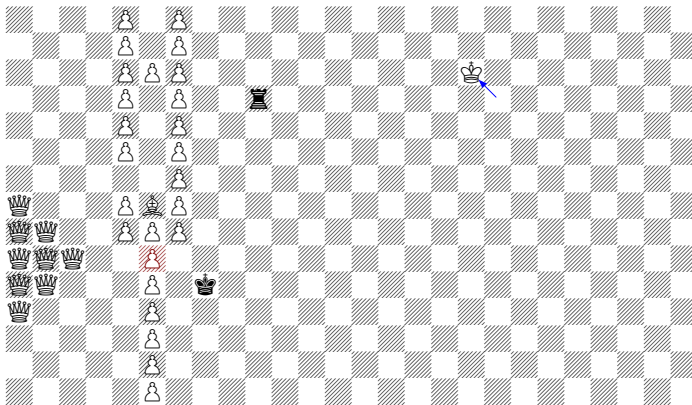
# Releasing the Hordes, with value $\omega^2$



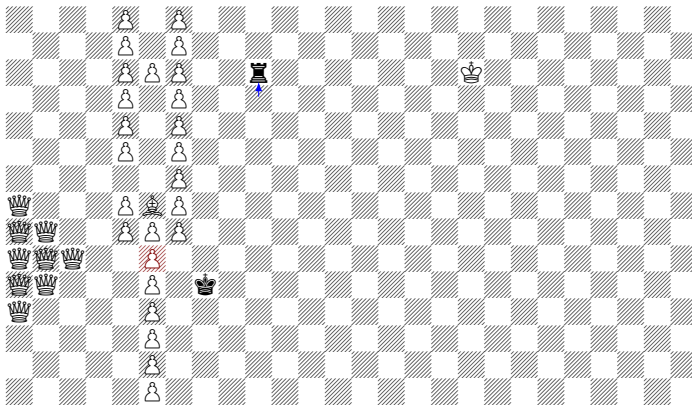




## Releasing the Hordes, with value $\omega^2$

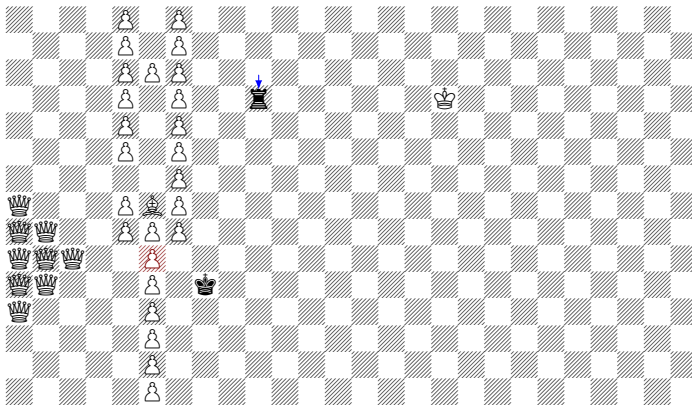


# Releasing the Hordes, with value $\omega^2$

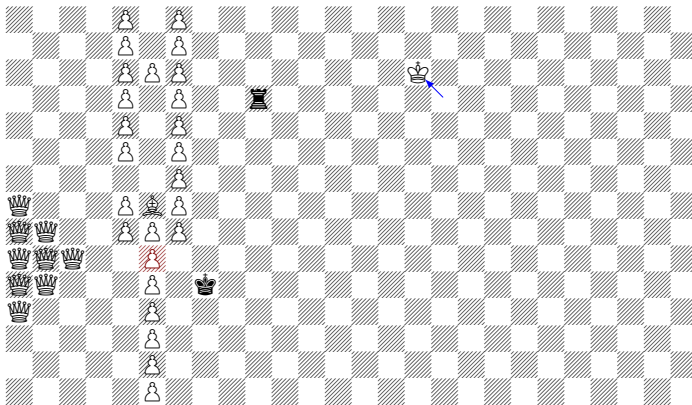




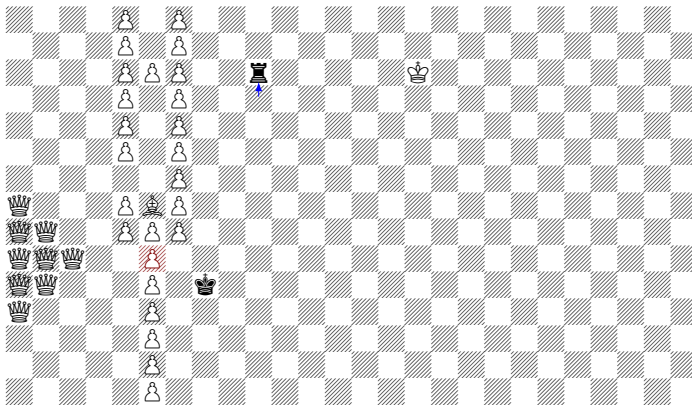
# Releasing the Hordes, with value $\omega^2$



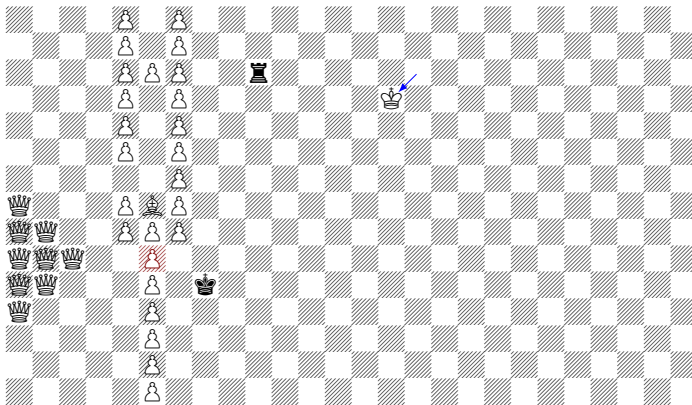
# Releasing the Hordes, with value $\omega^2$



## Releasing the Hordes, with value $\omega^2$

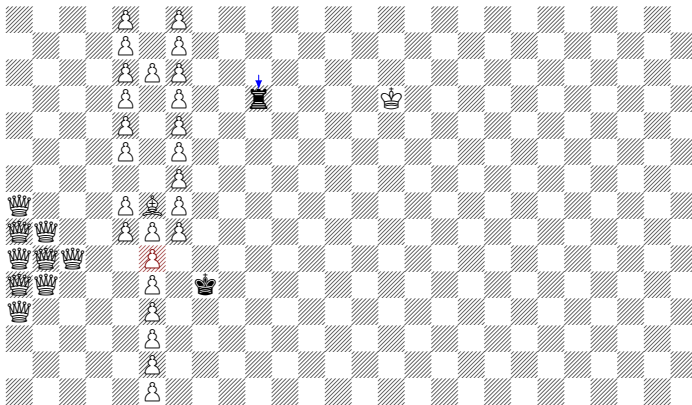


# Releasing the Hordes, with value $\omega^2$

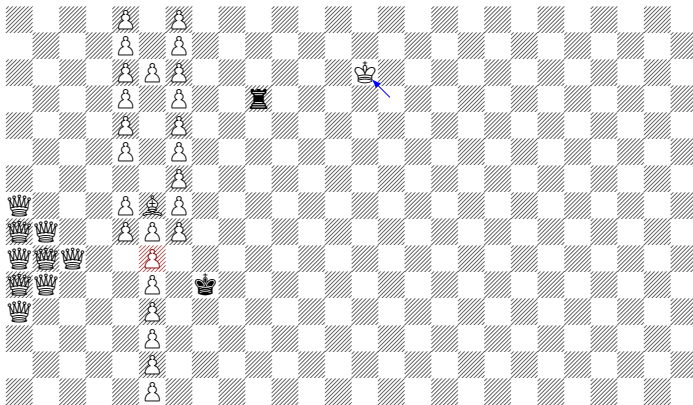




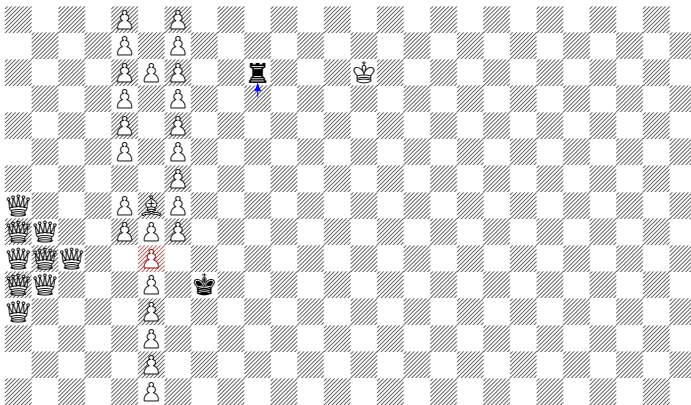
# Releasing the Hordes, with value $\omega^2$



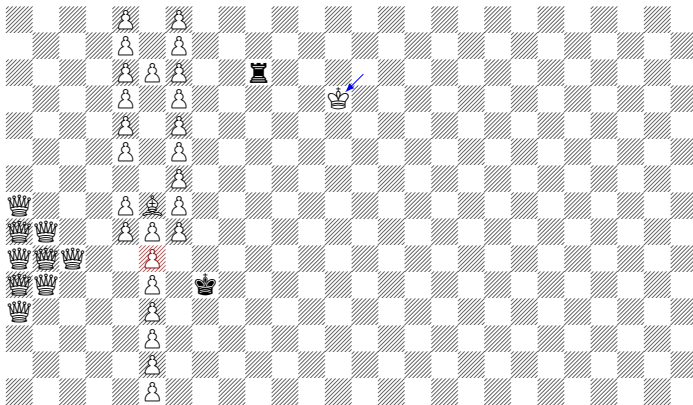
# Releasing the Hordes, with value $\omega^2$



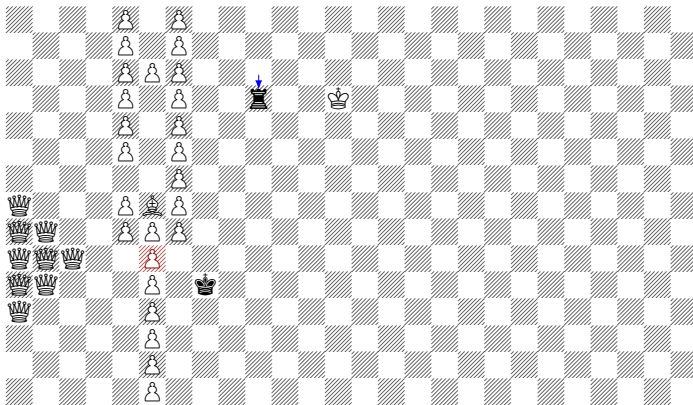
## Releasing the Hordes, with value $\omega^2$



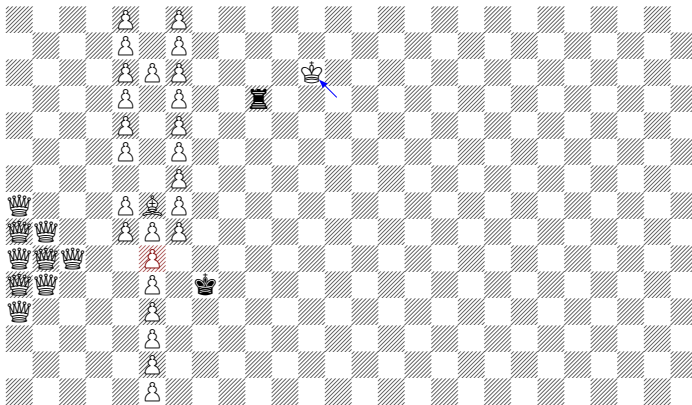
# Releasing the Hordes, with value $\omega^2$



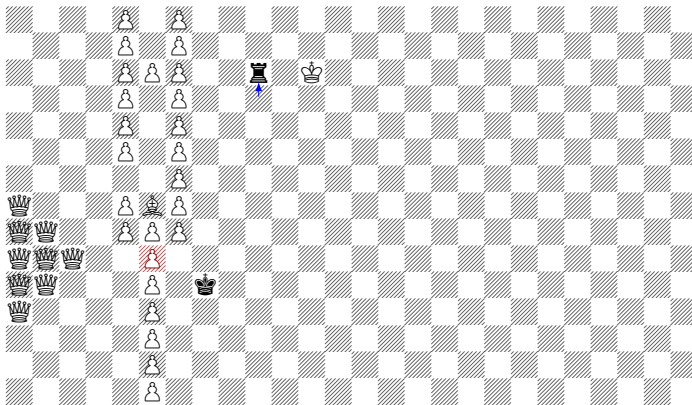
# Releasing the Hordes, with value $\omega^2$



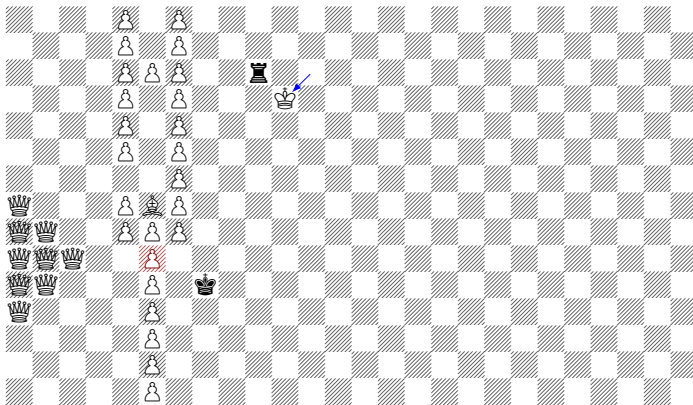
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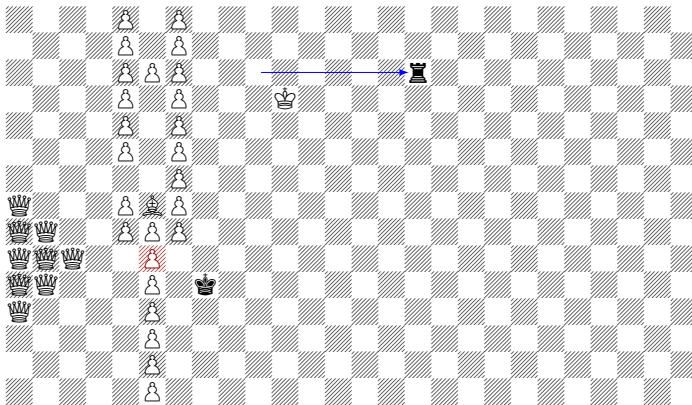


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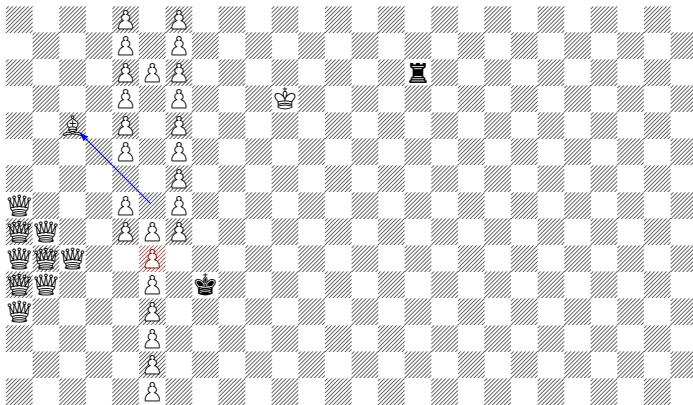




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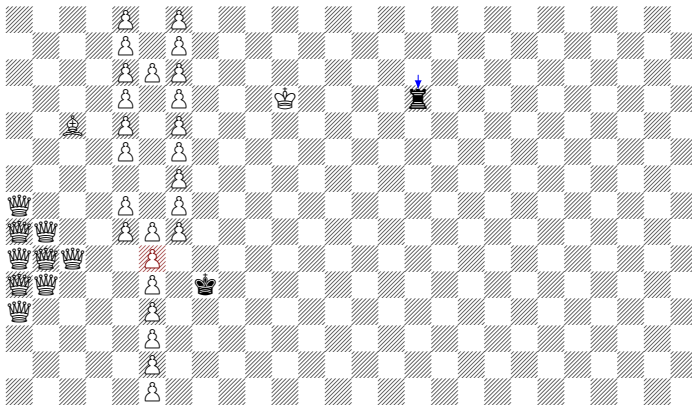


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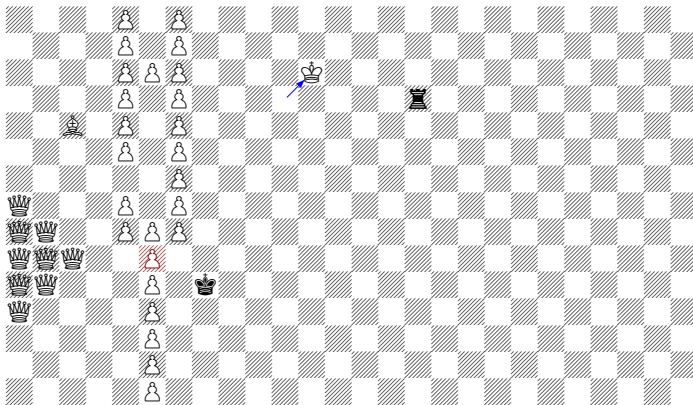


The bishop unlocks the door.

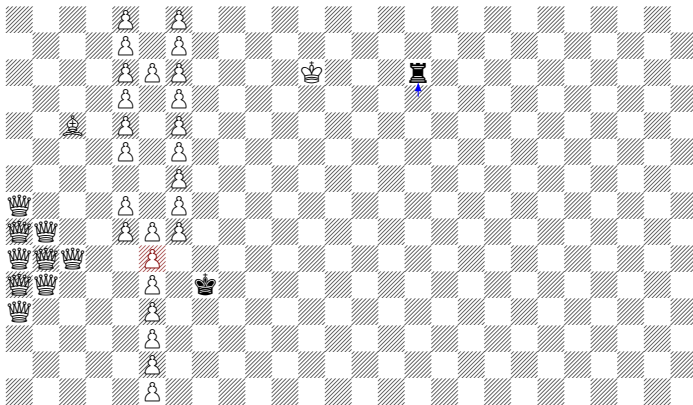
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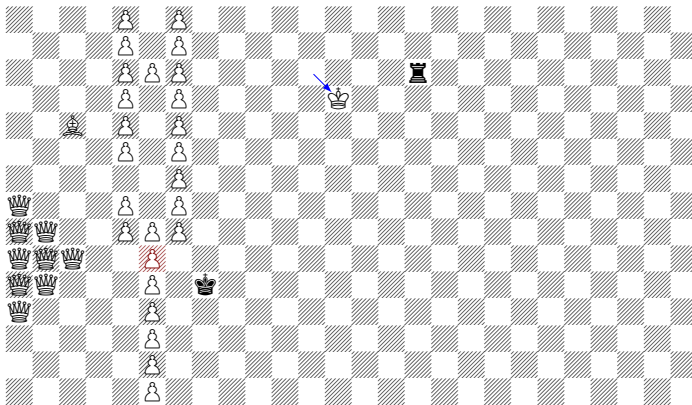
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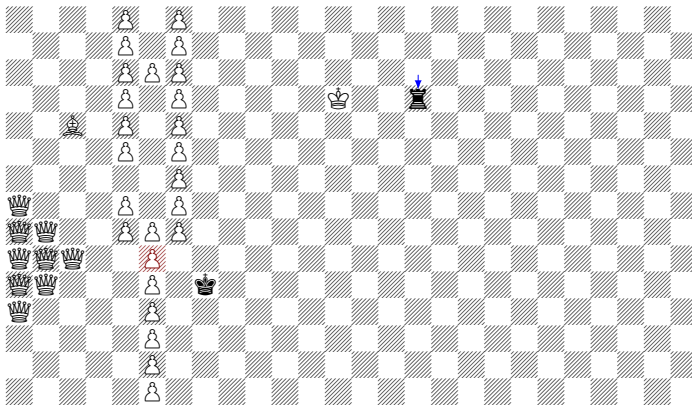
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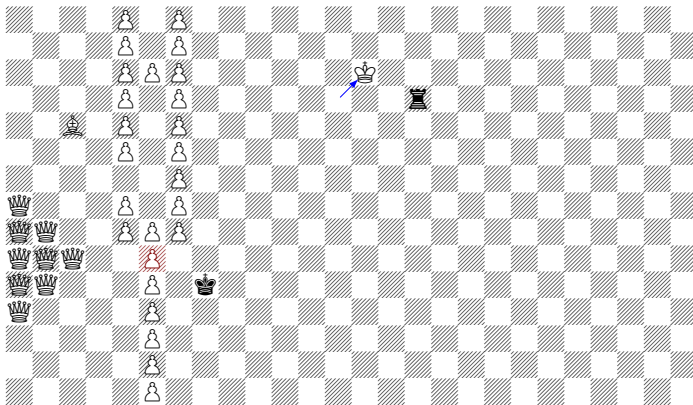
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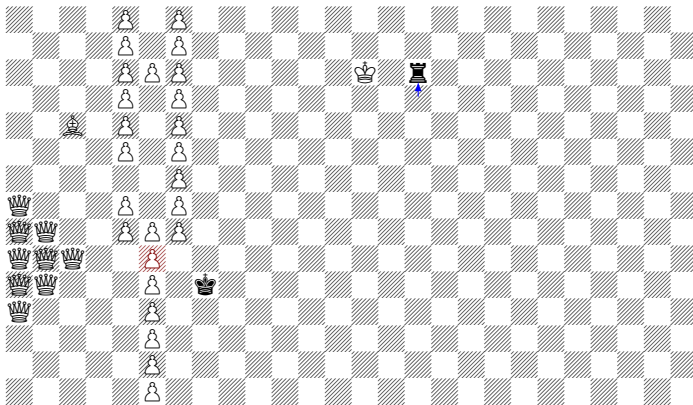


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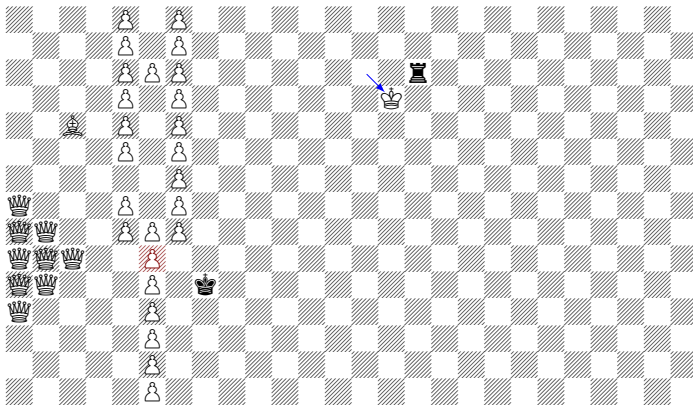




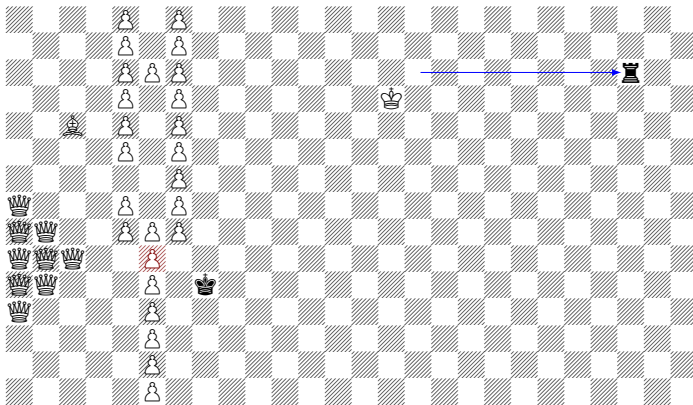
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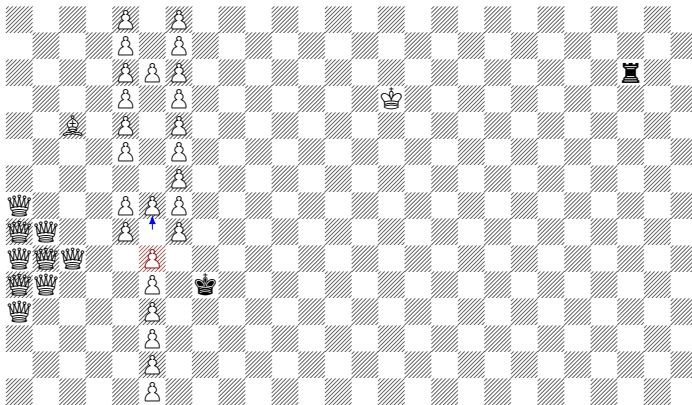


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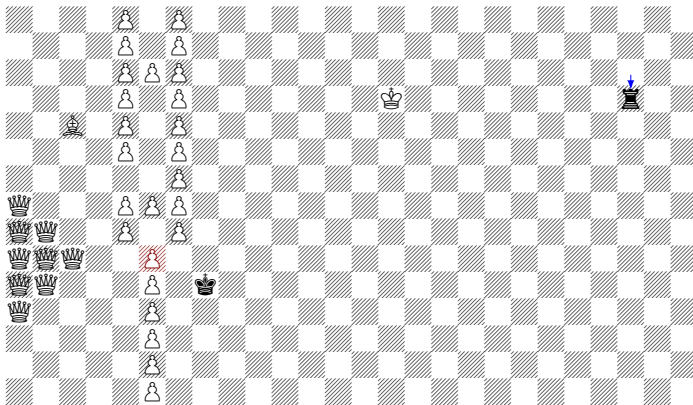


Black can move rook arbitrary distance.

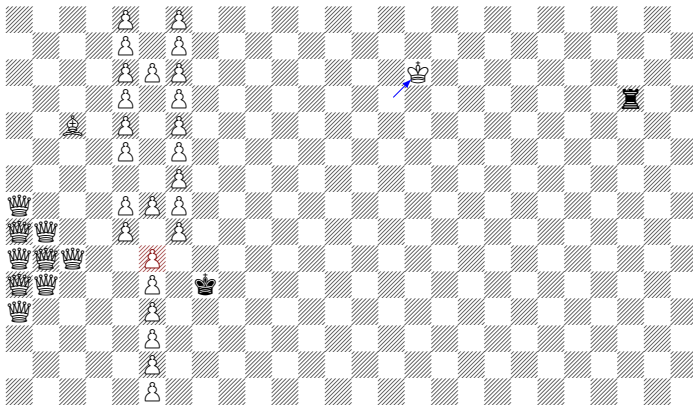
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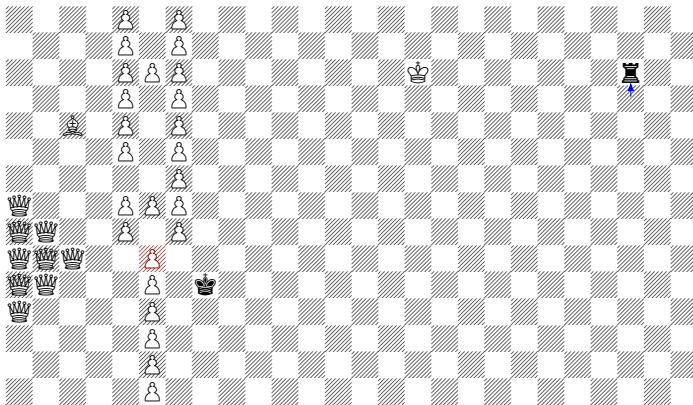
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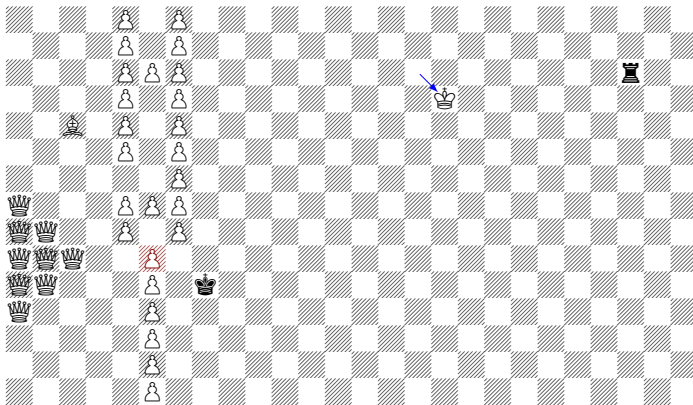
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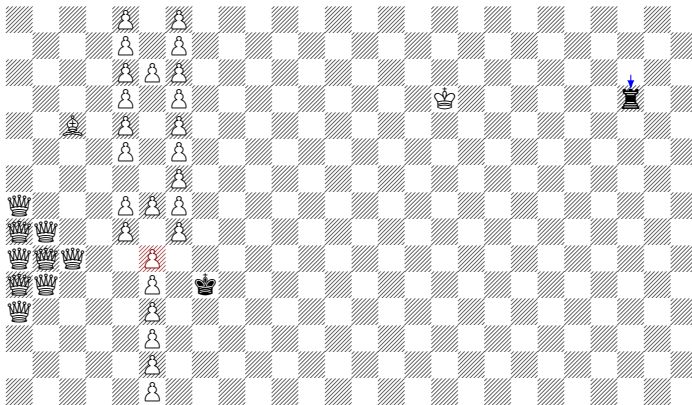


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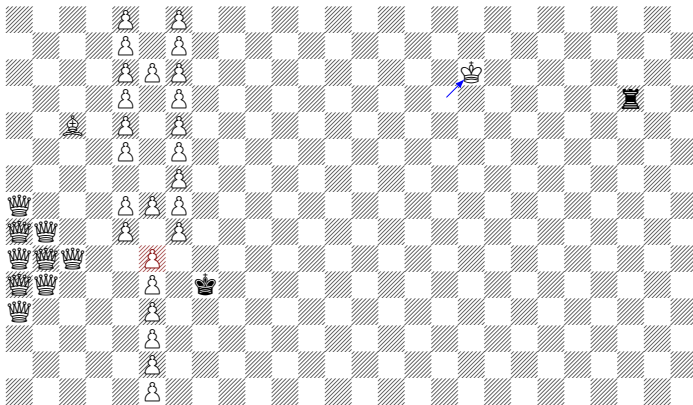




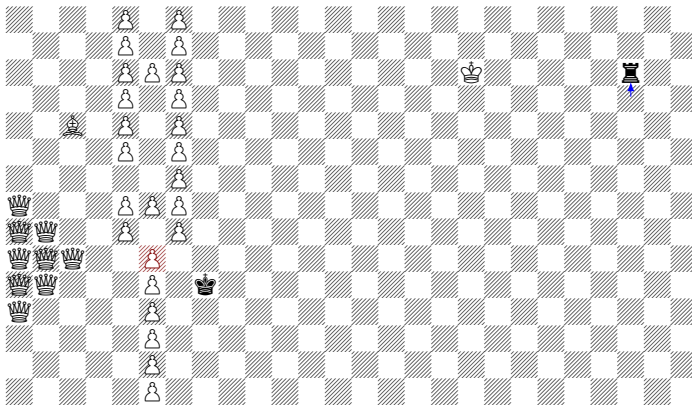
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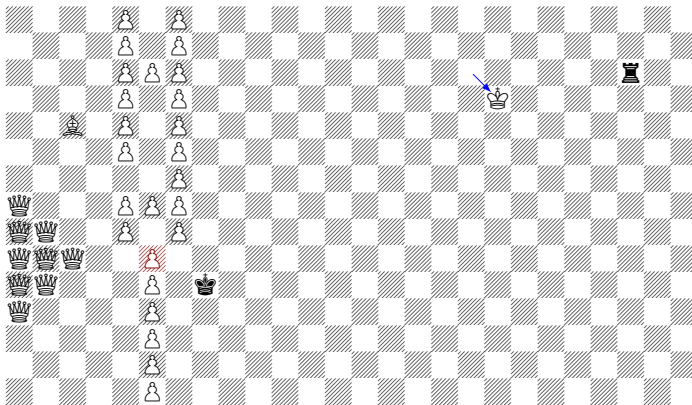
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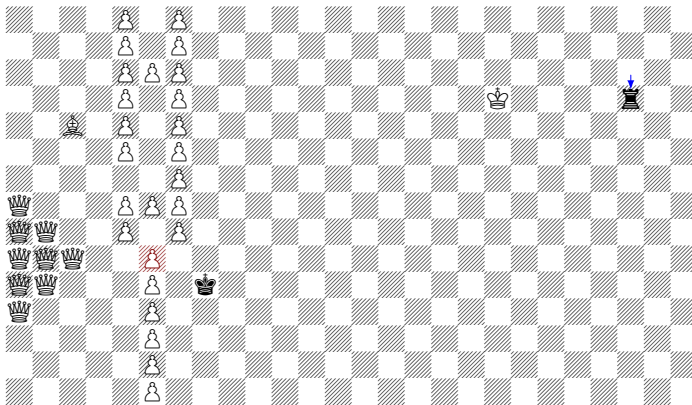
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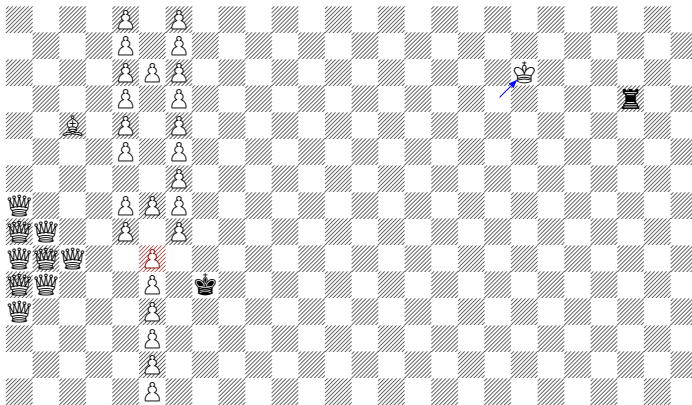
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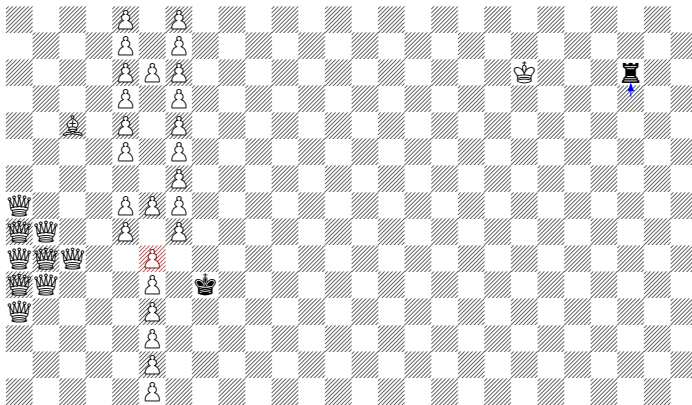
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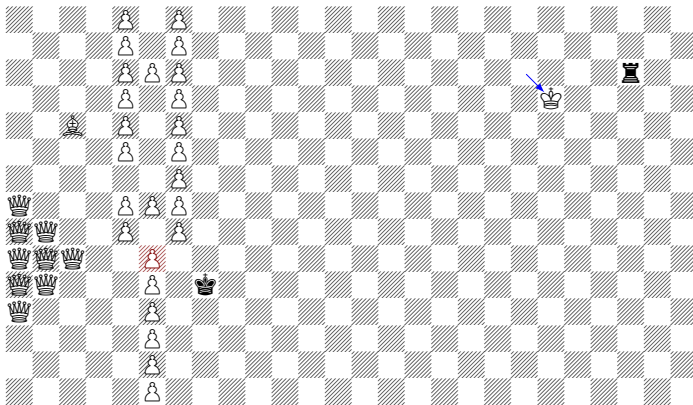
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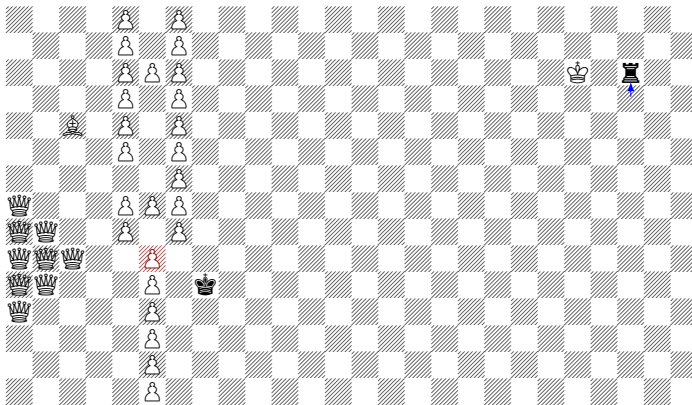
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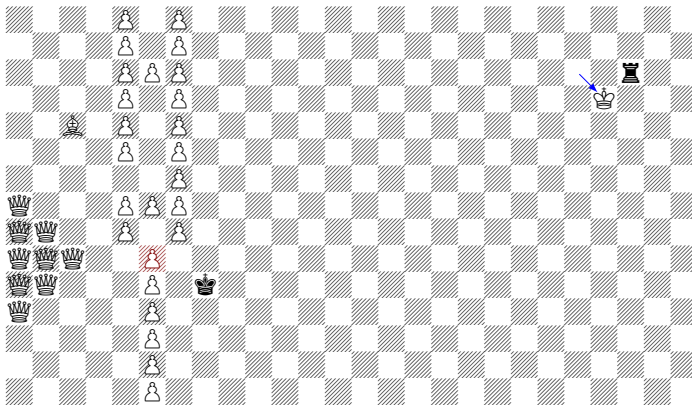
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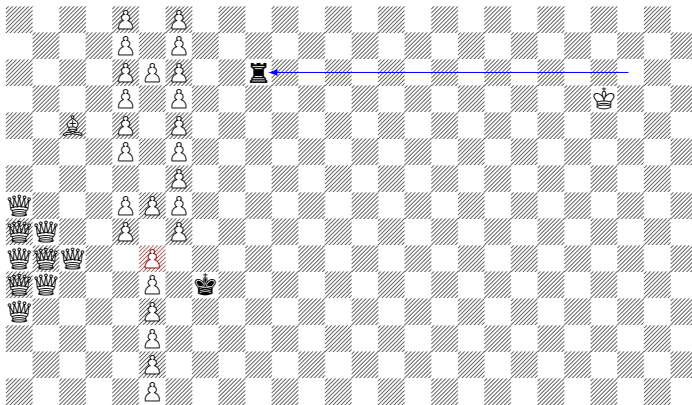
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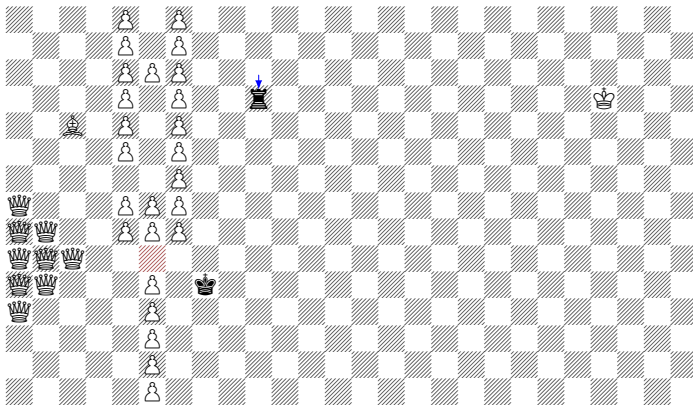


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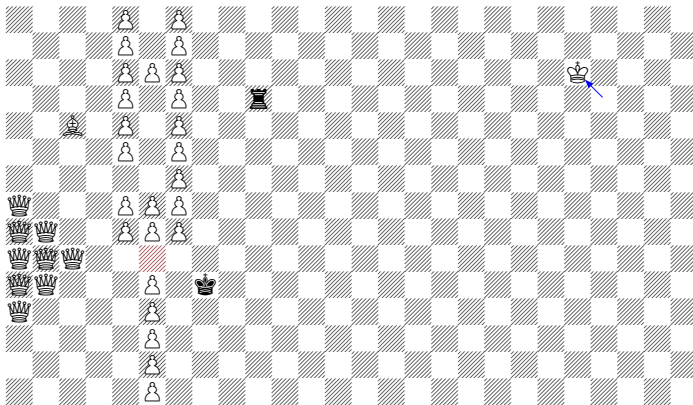


Joel David Hamkins

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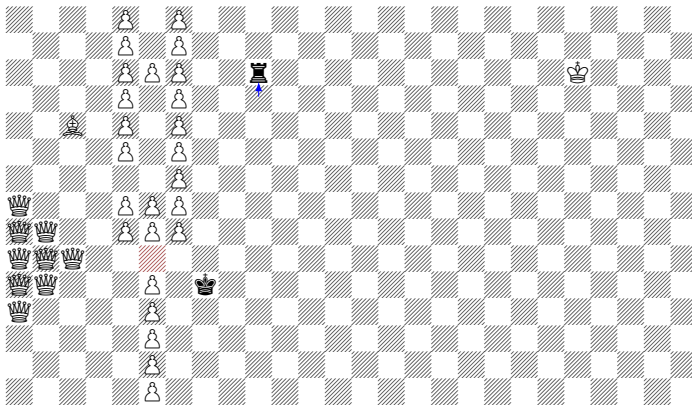


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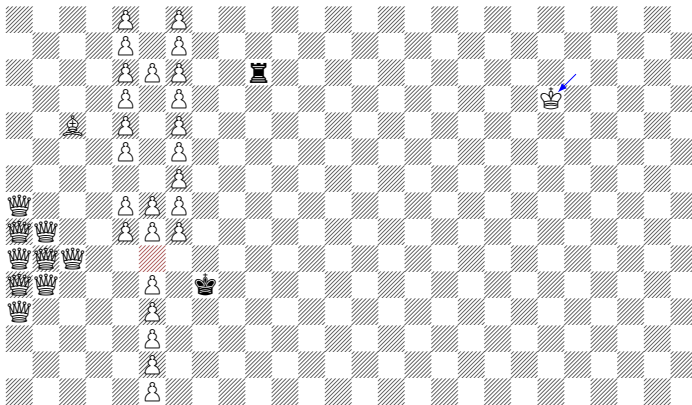




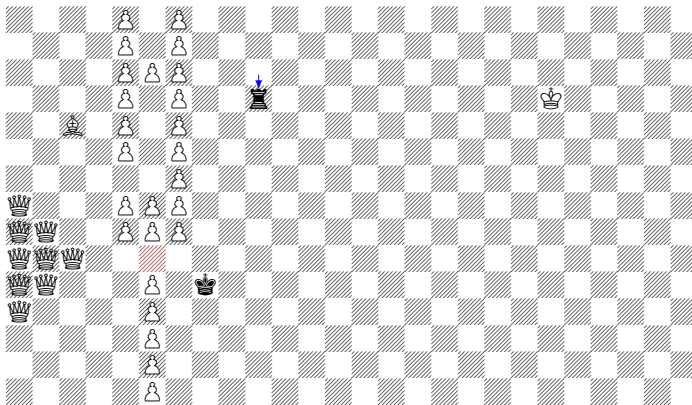
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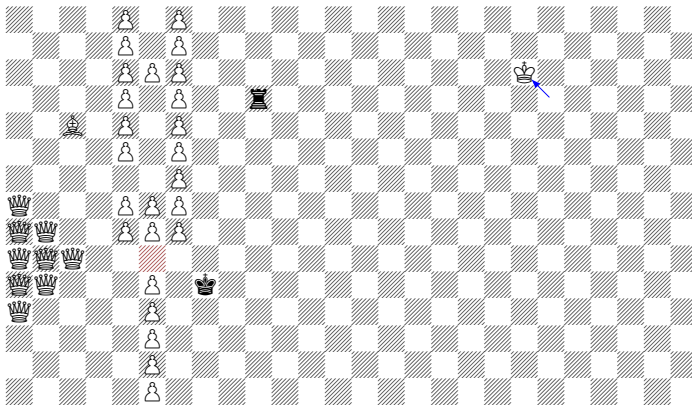
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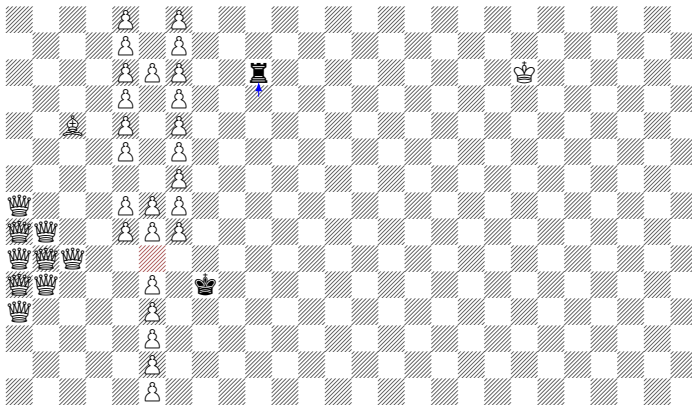
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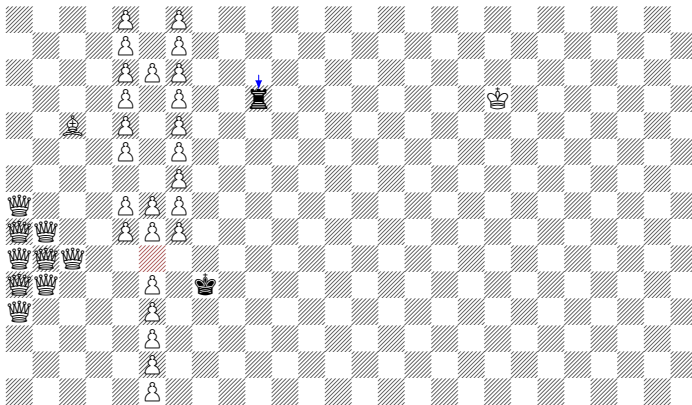


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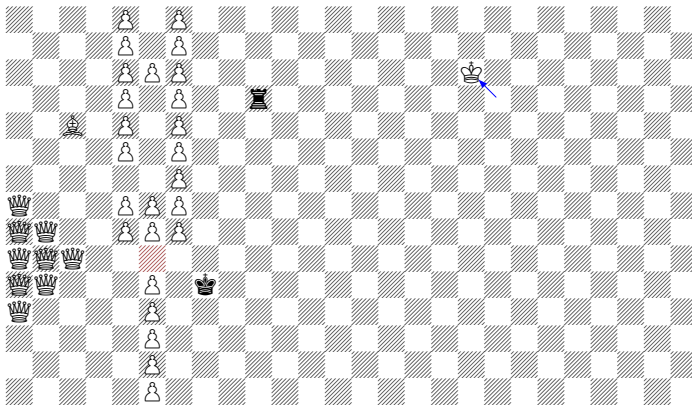


The chessboard diagram shows a complex position. White pieces are clustered on the left side, while Black pieces are scattered. A blue arrow points to the White King on g1.

# Releasing the Hordes, with value $\omega^2$

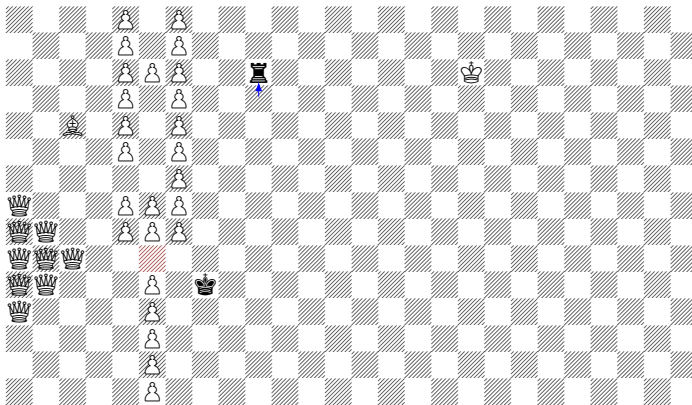


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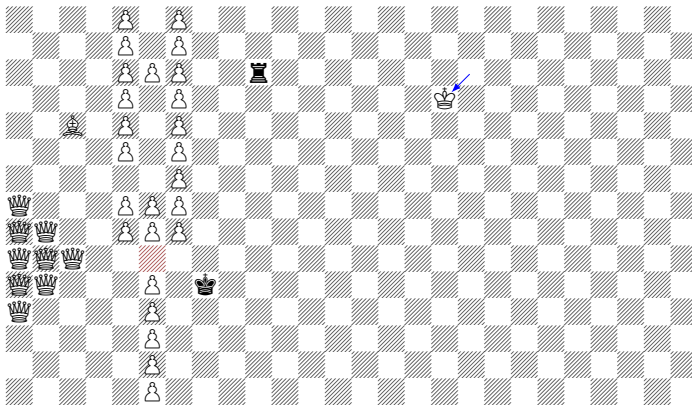




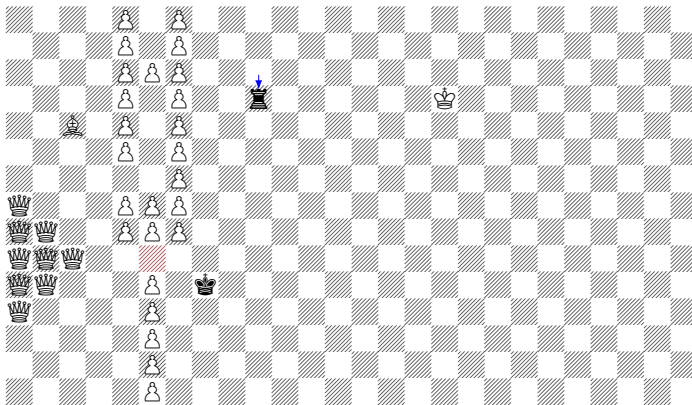
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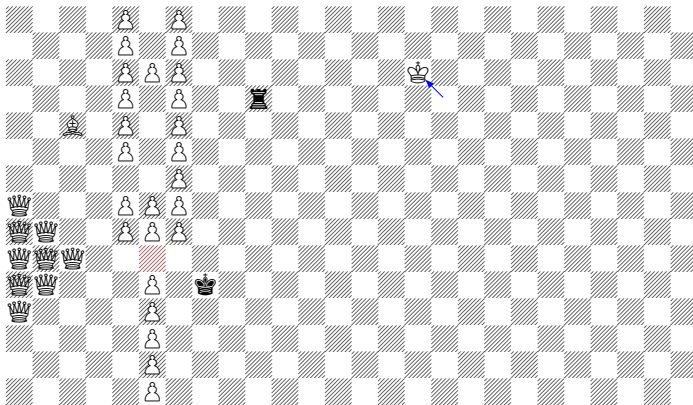
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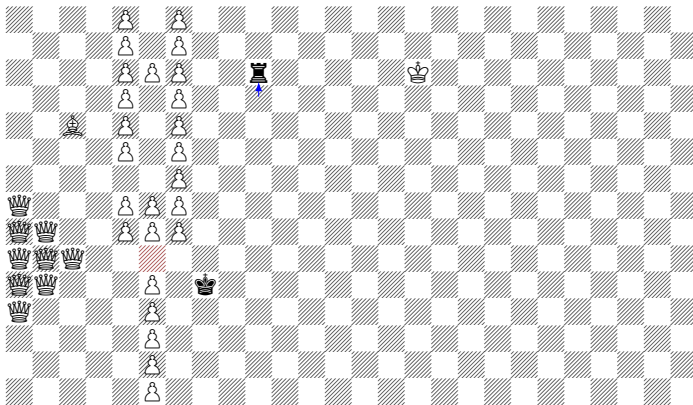
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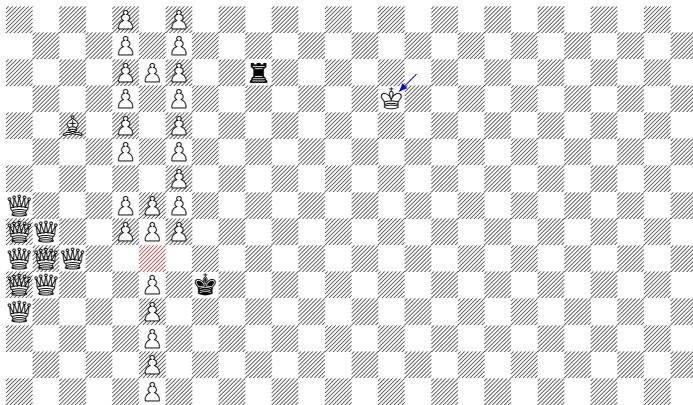
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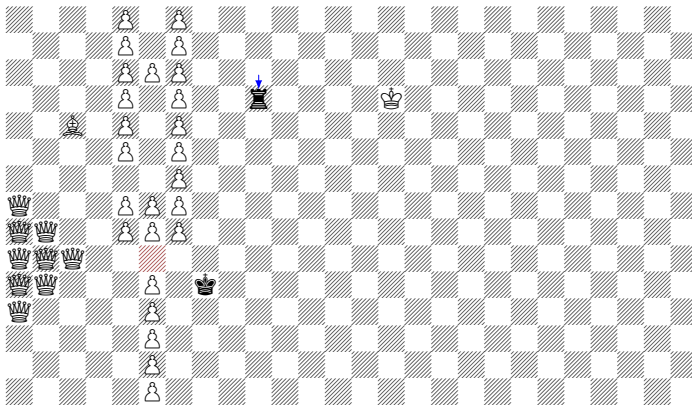
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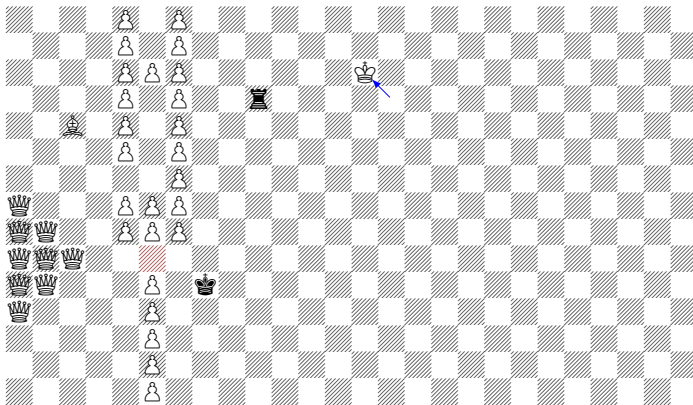
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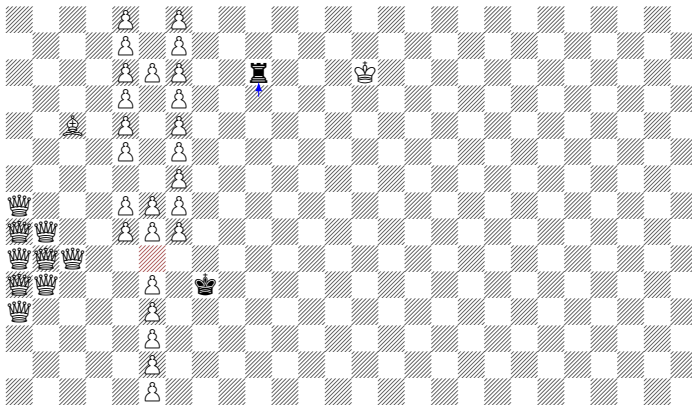


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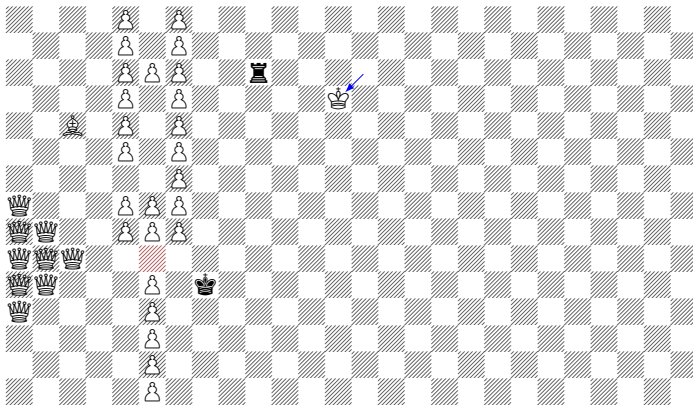




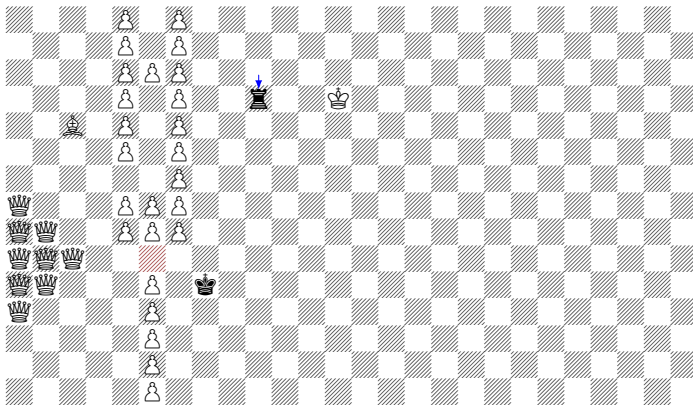
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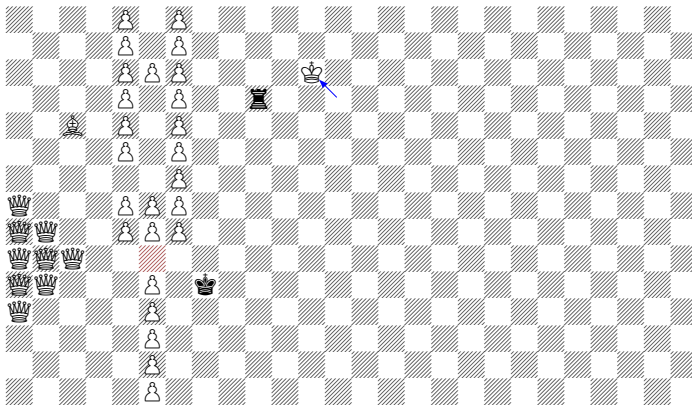
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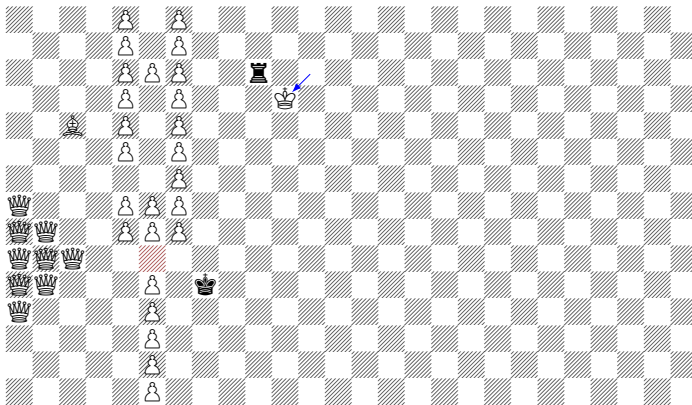


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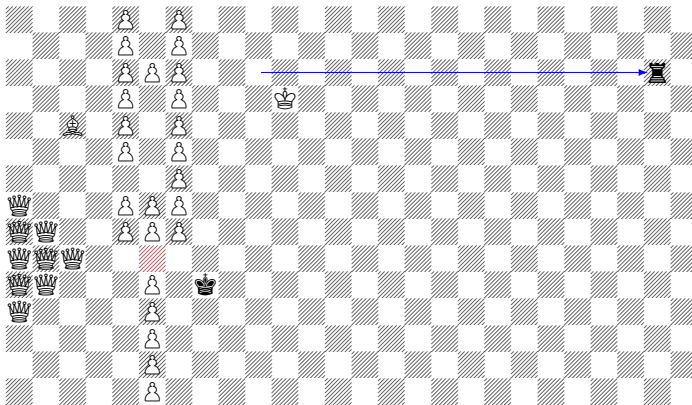




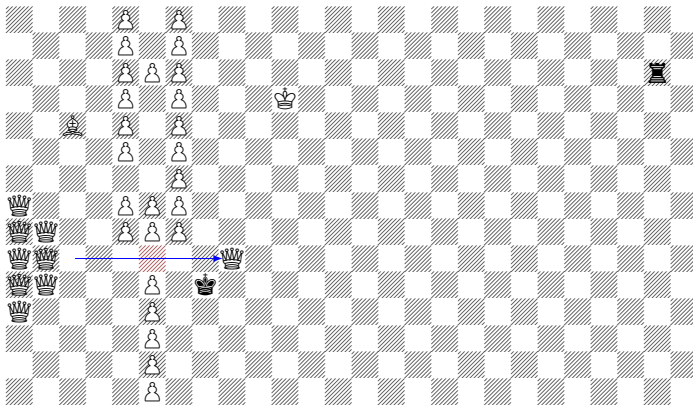
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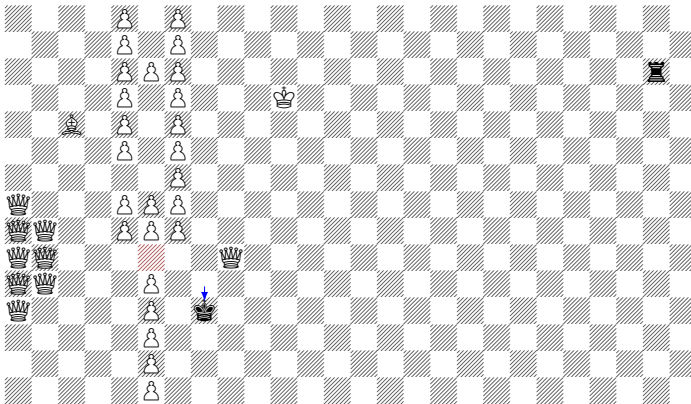
# Releasing the Hordes, with value $\omega^2$



Queens enter the mating chamber.

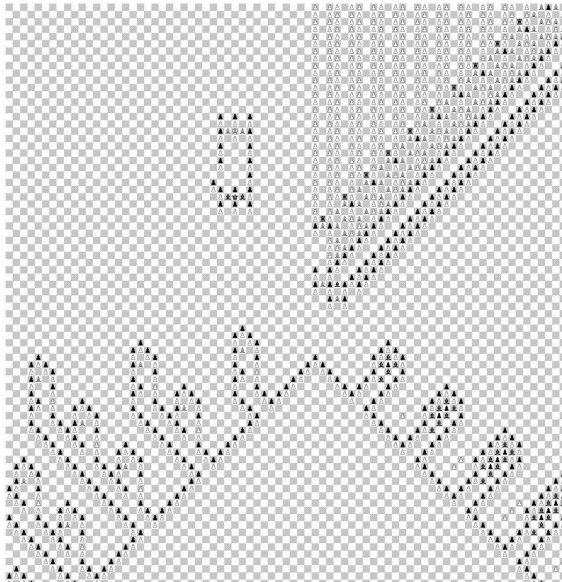


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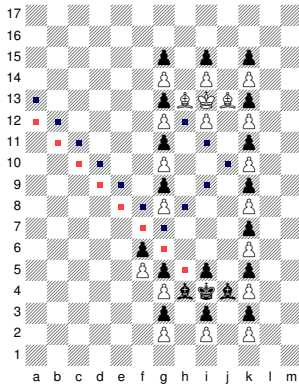




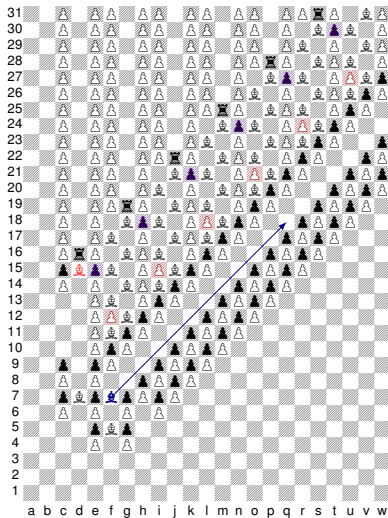
# State of the art: a position with value $\omega^4$



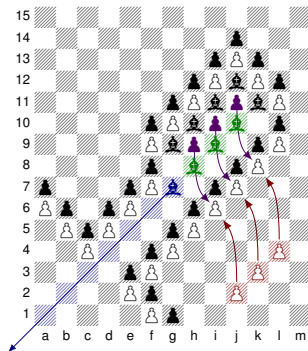
# The throne room



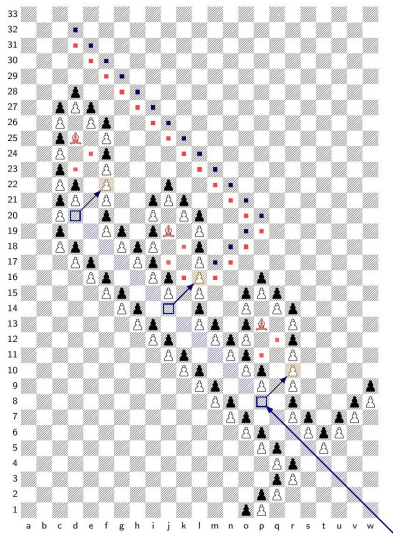
# The rook towers



# Bishop cannon



# Bishop gateway terminal



# 3D chess

Let us turn briefly to three-dimensional chess.



## 3D chess

The history of 3D chess spans more centuries than you expect:

- 19th century: Kieseritzky's Kubikschach in 1851.

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- 23rd century: Spock and Kirk play on several Star Trek episodes.



# Infinite 3D chess

We consider *infinite* three-dimensional chess, 3D chess on infinite boards with no boundary.

One must specify the piece movement; there is room for reasonable disagreement.

# Ordinal values in infinite 3D chess

Cory Evans and I proved that every countable ordinal arises as a game value in infinite 3D chess.

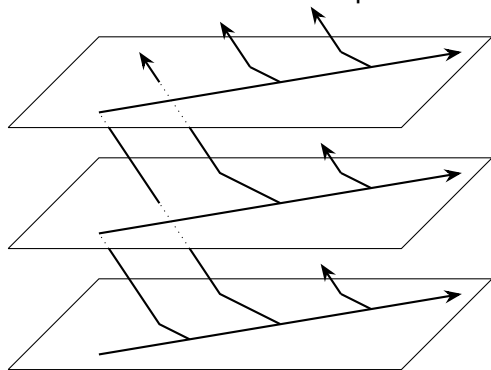
# Ordinal values in infinite 3D chess

Cory Evans and I proved that every countable ordinal arises as a game value in infinite 3D chess.

Let me discuss a few ideas from the proof.

# Embedding a tree into 3 space

Main idea is to embed complex trees into infinite 3D chess.



We embed the tree, with the infinite-branching nodes of  $T$  simulated individually on separate layers

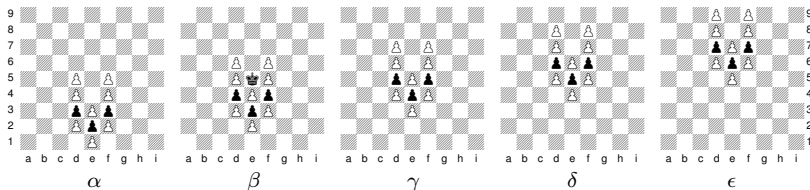


# Ascending stairs in 3D chess

Our positions involve embedded trees, with connections ascending to higher levels along various stairway passages.

# Ascending stairs in 3D chess

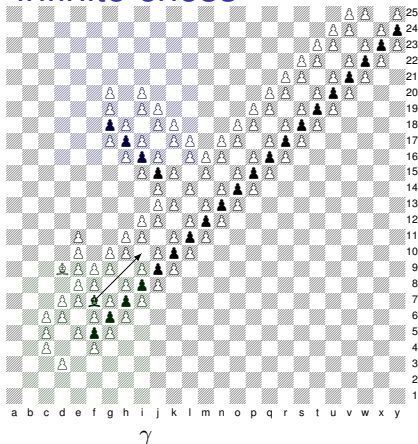
Our positions involve embedded trees, with connections ascending to higher levels along various stairway passages.



The black king is forced to ascend the stairs via

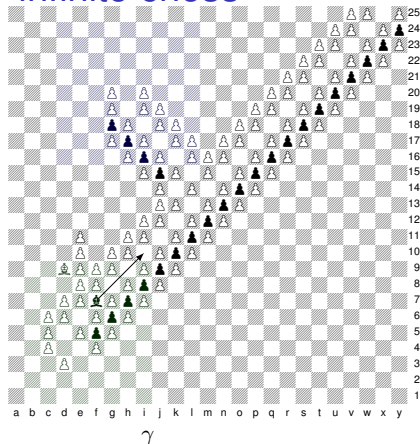
1.  $\alpha e4+$   $K\gamma e6$  2.  $\beta e5+$   $K\delta e7$  3.  $\gamma e6+$   $K\epsilon e8$  4.  $\delta e7+$ .

# Branching node layer in 3D infinite chess



Black king is forced to ascend stairway into this realm, where a choice will be made about which branching node to follow.

# Branching node layer in 3D infinite chess



Black king is forced to ascend stairway into this realm, where a choice will be made about which branching node to follow.

Main point: embedding trees  $\rightarrow$  high game values

Thank you.



Slides and articles available on <http://jdh.hamkins.org>.

Joel David Hamkins  
University of Notre Dame

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